

Bit Arithmetics

CS143: lecture 15

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Unsigned Arithmetics

- unsigned means without sign (+/-). In C, `unsigned int x = 4`
- Why unsigned?
 - 32 bits means 2^{32} choices
 - If we don't allow negative number, we have range $[0, 2^{32} - 1]$
 - If we have negative numbers, the maximum has to be smaller.
- Most things that we care about are non-negative
 - counts, lengths, indices, dimensions...

Unsigned Arithmetics

Operators

- Arithmetic operators: $+$, $-$, $*$, $/$
- Bitwise operators: $\&$, $|$, $\wedge$, $\sim$
- Shifts: $\ll$, $\gg$

Unsigned Arithmetics

```
unsigned int n = 45;
```

n	2^n
0	1
1	2
2	4
3	8
4	16
5	32
6	64
7	128

1101 0011 1100 0001 0010 0100 1001 1111

& 0101 0101 0101 0101 0101 0101 0101 0101

Unsigned Arithmetics

```
unsigned int n = 45;
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& 0101 0101 0101 0101 0101 0101 0101 0101

1

Unsigned Arithmetics

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& 0101 0101 0101 0101 0101 0101 0101 0101

11

Unsigned Arithmetics

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111

Unsigned Arithmetics

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1101 0011 1100 0001 0010 0100 1001

& 0101 0101 0101 0101 0101 0101 0101

0111

0111

```
unsigned int n = 45;
```

n	2^n
0	1
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3	8
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7	128

	1101	0011	1100	0001	0010	0100	1001	1111
&	0101	0101	0101	0101	0101	0101	0101	0101
								1 0111

```
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	1101	0011	1100	0001	0010	0100	1001	1111
&	0101	0101	0101	0101	0101	0101	0101	0101
							01	0111

Unsigned Arithmetics

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```

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7	128

1101 0011 1100 0001 0010 0100 1001 1111

& 0101 0101 0101 0101 0101 0101 0101 0101

0101 0001 0100 0001 0000 0100 0001 0101

Unsigned Arithmetics

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1101 0011 1100 0001 0010 0100 1001 1111

| 0101 0101 0101 0101 0101 0101 0101 0101

1

Unsigned Arithmetics

```
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7	128

1101 0011 1100 0001 0010 0100 1001 1111

| 0101 0101 0101 0101 0101 0101 0101 0101

11

Unsigned Arithmetics

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| 0101 0101 0101 0101 0101 0101 0101 0101

111

Unsigned Arithmetics

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6	64
7	128

1101 0011 1100 0001 0010 0100 1001

| 0101 0101 0101 0101 0101 0101 0101

1111

1111

```
unsigned int n = 45;
```

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unsigned int n = 45;
```

n	2^n
0	1
1	2
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	1101	0011	1100	0001	0010	0100	1001	1111
	0101	0101	0101	0101	0101	0101	0101	0101
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	1101	0011	1100	0001	0010	0100	1001	1111
	0101	0101	0101	0101	0101	0101	0101	0101
							01	1111

Unsigned Arithmetics

```
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n	2^n
0	1
1	2
2	4
3	8
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5	32
6	64
7	128

1101 0011 1100 0001 0010 0100 1001 1111

| 0101 0101 0101 0101 0101 0101 0101 0101

101 1111

Unsigned Arithmetics

```
unsigned int n = 45;
```

n	2^n
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1	2
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7	128

1101 0011 1100 0001 0010 0100 1001 1111

| 0101 0101 0101 0101 0101 0101 0101 0101

1101 0111 1101 0101 0111 0101 1101 1111

Unsigned Arithmetics

```
unsigned int n = 45;
```

n	2^n
0	1
1	2
2	4
3	8
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5	32
6	64
7	128

~ 1101 0011 1100 0001 0010 0100 1001 1111

Unsigned Arithmetics

```
unsigned int n = 45;
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n	2^n
0	1
1	2
2	4
3	8
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5	32
6	64
7	128

~ 1101 0011 1100 0001 0010 0100 1001 1111

0000

Unsigned Arithmetics

```
unsigned int n = 45;
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n	2^n
0	1
1	2
2	4
3	8
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5	32
6	64
7	128

~ 1101 0011 1100 0001 0010 0100 1001 1111

0110 0000

Unsigned Arithmetics

```
unsigned int n = 45;
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n	2^n
0	1
1	2
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6	64
7	128

~ 1101 0011 1100 0001 0010 0100 1001 1111

1011 0110 0000

Unsigned Arithmetics

```
unsigned int n = 45;
```

n	2^n
0	1
1	2
2	4
3	8
4	16
5	32
6	64
7	128

~ 1101 0011 1100 0001 0010 0100 1001 1111

0010 1100 0011 1110 1101 1011 0110 0000

Unsigned Arithmetics

```
unsigned int n = 45;
```

n	2^n
0	1
1	2
2	4
3	8
4	16
5	32
6	64
7	128

Exclusive-Or (XOR): 1 if the bits are different
0 if the bits are the same.

1101 0011 1100 0001 0010 0100 1001 1111

^

0101 0101 0101 0101 0101 0101 0101 0101

Unsigned Arithmetics

```
unsigned int n = 45;
```

n	2^n
0	1
1	2
2	4
3	8
4	16
5	32
6	64
7	128

1101 0011 1100 0001 0010 0100 1001 1111

^ 0101 0101 0101 0101 0101 0101 0101 0101

0

```
unsigned int n = 45;
```

n	2^n
0	1
1	2
2	4
3	8
4	16
5	32
6	64
7	128

[illegible]

```
unsigned int n = 45;
```

n	2^n
0	1
1	2
2	4
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5	32
6	64
7	128

	1101	0011	1100	0001	0010	0100	1001	1111
^	0101	0101	0101	0101	0101	0101	0101	0101
								010

Unsigned Arithmetics

```
unsigned int n = 45;
```

n	2^n
0	1
1	2
2	4
3	8
4	16
5	32
6	64
7	128

1101 0011 1100 0001 0010 0100 1001 1111

^ 0101 0101 0101 0101 0101 0101 0101 0101

1010

Unsigned Arithmetics

```
unsigned int n = 45;
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n	2^n
0	1
1	2
2	4
3	8
4	16
5	32
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7	128

1101 0011 1100 0001 0010 0100 1001 1111

^

0101 0101 0101 0101 0101 0101 0101 0101

0 1010

Unsigned Arithmetics

```
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n	2^n
0	1
1	2
2	4
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4	16
5	32
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7	128

1101 0011 1100 0001 0010 0100 1001 1111

^

0101 0101 0101 0101 0101 0101 0101 0101

10 1010

Unsigned Arithmetics

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n	2^n
0	1
1	2
2	4
3	8
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5	32
6	64
7	128

1101 0011 1100 0001 0010 0100 1001 1111

^

0101 0101 0101 0101 0101 0101 0101 0101

100 1010

Unsigned Arithmetics

```
unsigned int n = 45;
```

n	2^n
0	1
1	2
2	4
3	8
4	16
5	32
6	64
7	128

1101 0011 1100 0001 0010 0100 1001 1111

^

0101 0101 0101 0101 0101 0101 0101 0101

1000 0110 1001 0100 0111 0001 1100 1010

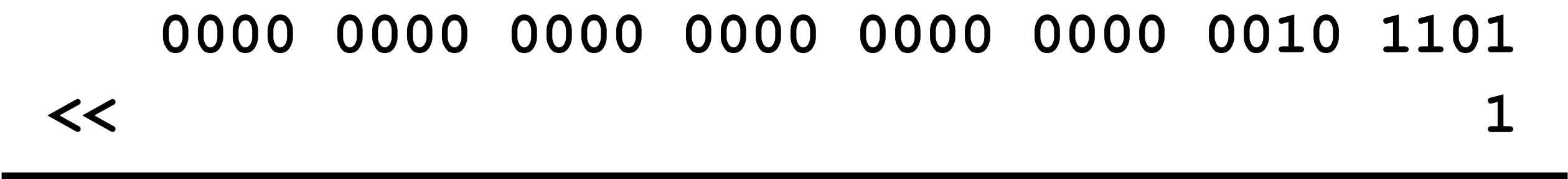
Unsigned Arithmetics

- $\sim$ is not the same as $!$, $\&$ is not the same as $\&\&$, $|$ is not the same as $||$
 - $!$, $\&\&$, $||$ are *Boolean* operators -- they treat zero as false and non-zero as true
 - $\sim$, $\&$, $|$ are *bitwise* operators -- they operate on each bit.

Unsigned Arithmetics

```
unsigned int n = 45;
```

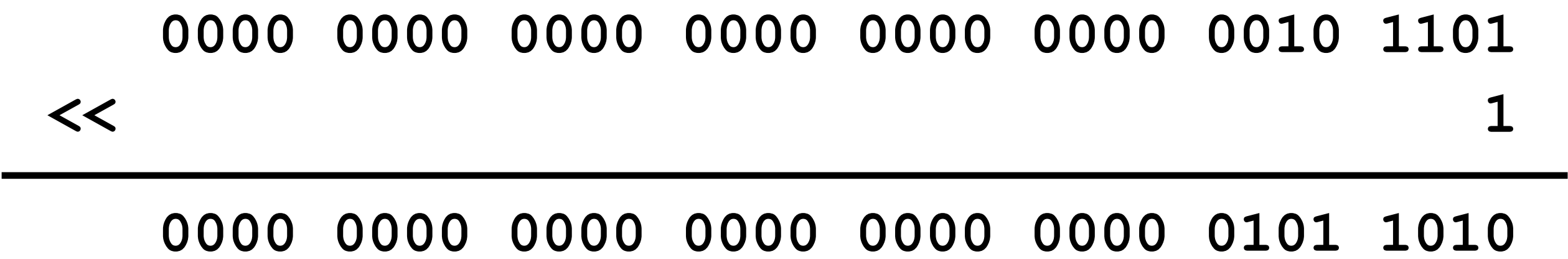
n	2^n
0	1
1	2
2	4
3	8
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5	32
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Unsigned Arithmetics

```
unsigned int n = 45;
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n	2^n
0	1
1	2
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Unsigned Arithmetics

```
unsigned int n = 45;
```

n	2^n
0	1
1	2
2	4
3	8
4	16
5	32
6	64
7	128

0000 0000 0000 0000 0000 0000 0010 1101

<<1

0000 0000 0000 0000 0000 0000 0101 1010

= 0x5A

Unsigned Arithmetics

```
unsigned int n = 45;
```

n	2^n
0	1
1	2
2	4
3	8
4	16
5	32
6	64
7	128

0000 0000 0000 0000 0000 0000 0010 1101

<<1

0000 0000 0000 0000 0000 0000 0101 1010

= 0x5A = 90

Unsigned Arithmetics

```
unsigned int n = 45;
```

n	2^n
0	1
1	2
2	4
3	8
4	16
5	32
6	64
7	128

0000 0000 0000 0000 0000 0000 0010 1101

<<1

0000 0000 0000 0000 0000 0000 0101 1010

<<16

= 0x5A = 90

Unsigned Arithmetics

```
unsigned int n = 45;
```

n	2^n
0	1
1	2
2	4
3	8
4	16
5	32
6	64
7	128

0000 0000 0000 0000 0000 0000 0010 1101

<<

1

0000 0000 0000 0000 0000 0000 0101 1010

<<

16

0000 0000 0101 1010 0000 0000 0000 0000

= 0x5A = 90

Unsigned Arithmetics

```
unsigned int n = 45;
```

n	2^n
0	1
1	2
2	4
3	8
4	16
5	32
6	64
7	128

0000 0000 0000 0000 0000 0000 0010 1101

<<

1

0000 0000 0000 0000 0000 0000 0101 1010

<<

16

0000 0000 0101 1010 0000 0000 0000 0000

<<

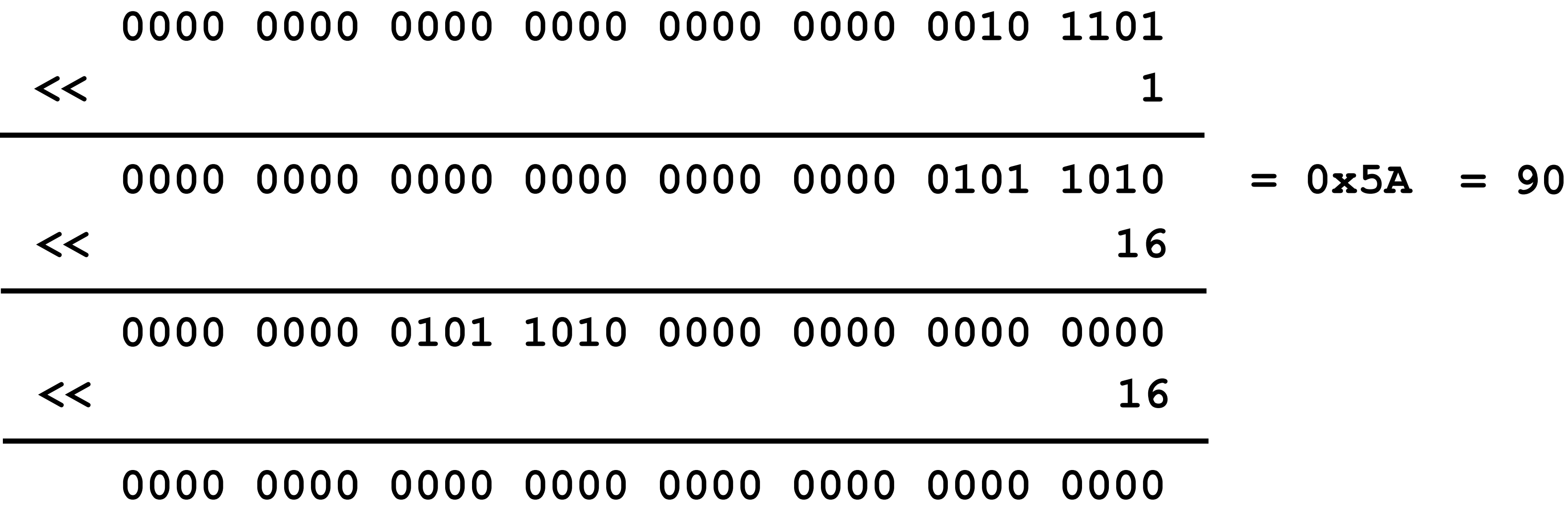
16

= 0x5A = 90

Unsigned Arithmetics

```
unsigned int n = 45;
```

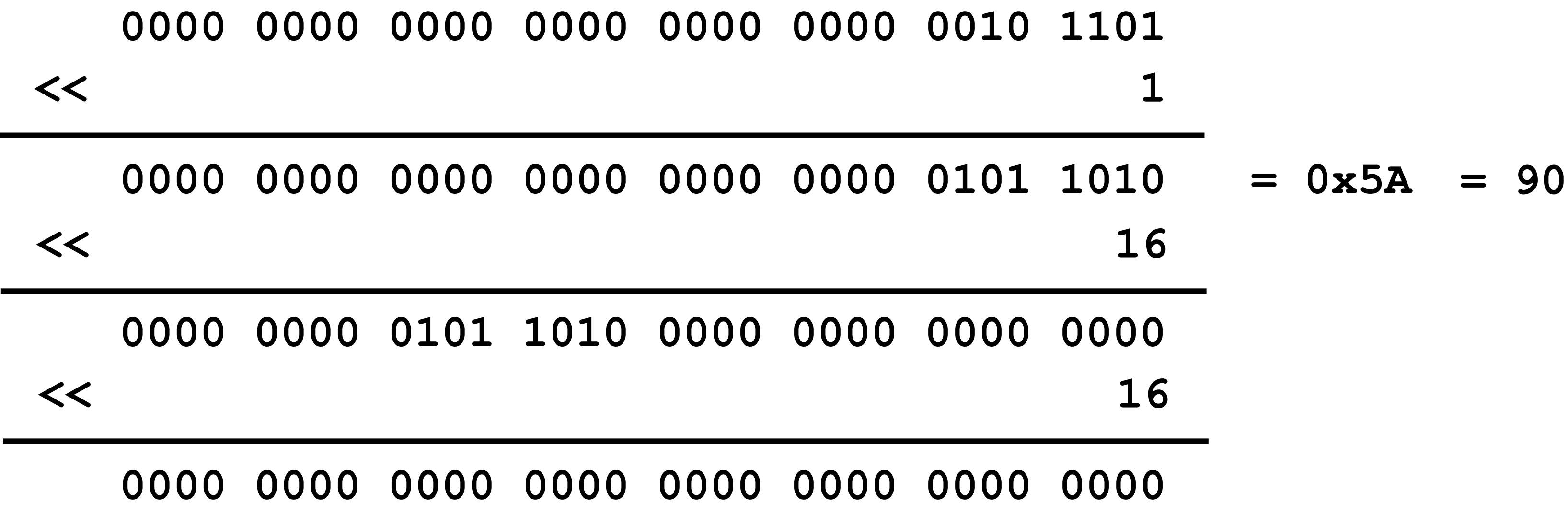
n	2^n
0	1
1	2
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Unsigned Arithmetics

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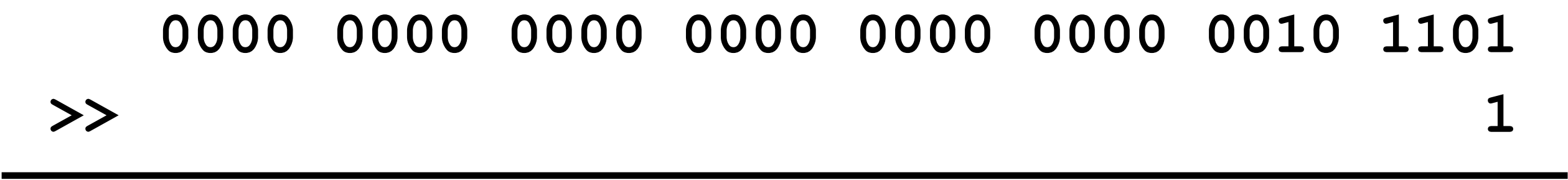


- When shift overflows, the bit is lost.

Unsigned Arithmetics

```
unsigned int n = 45;
```

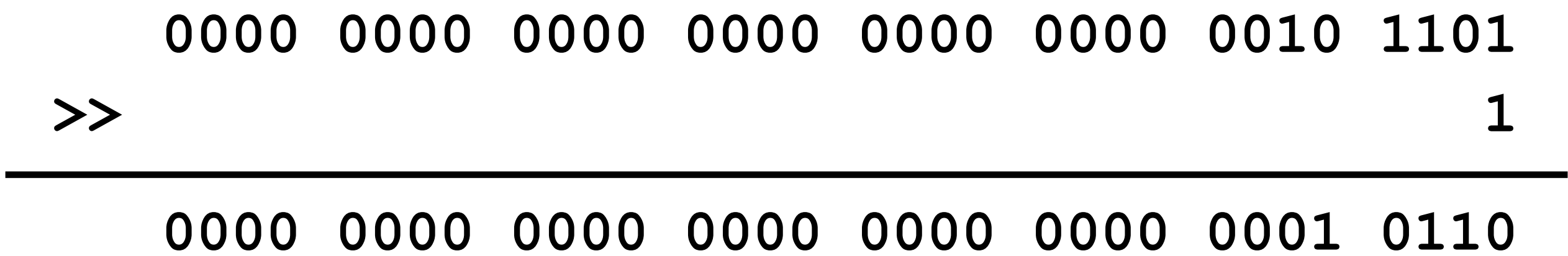
n	2^n
0	1
1	2
2	4
3	8
4	16
5	32
6	64
7	128



Unsigned Arithmetics

```
unsigned int n = 45;
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n	2^n
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2	4
3	8
4	16
5	32
6	64
7	128



Unsigned Arithmetics

```
unsigned int n = 45;
```

n	2^n
0	1
1	2
2	4
3	8
4	16
5	32
6	64
7	128

0000 0000 0000 0000 0000 0000 0010 1101

>> 1

0000 0000 0000 0000 0000 0000 0001 0110

= 0x16

Unsigned Arithmetics

```
unsigned int n = 45;
```

n	2^n
0	1
1	2
2	4
3	8
4	16
5	32
6	64
7	128

0000 0000 0000 0000 0000 0000 0010 1101

>> 1

0000 0000 0000 0000 0000 0000 0001 0110

= 0x16 = 22

Unsigned Arithmetics

```
unsigned int n = 45;
```

n	2^n
0	1
1	2
2	4
3	8
4	16
5	32
6	64
7	128

0000 0000 0000 0000 0000 0000 0010 1101

>> 1

0000 0000 0000 0000 0000 0000 0001 0110

>> 8

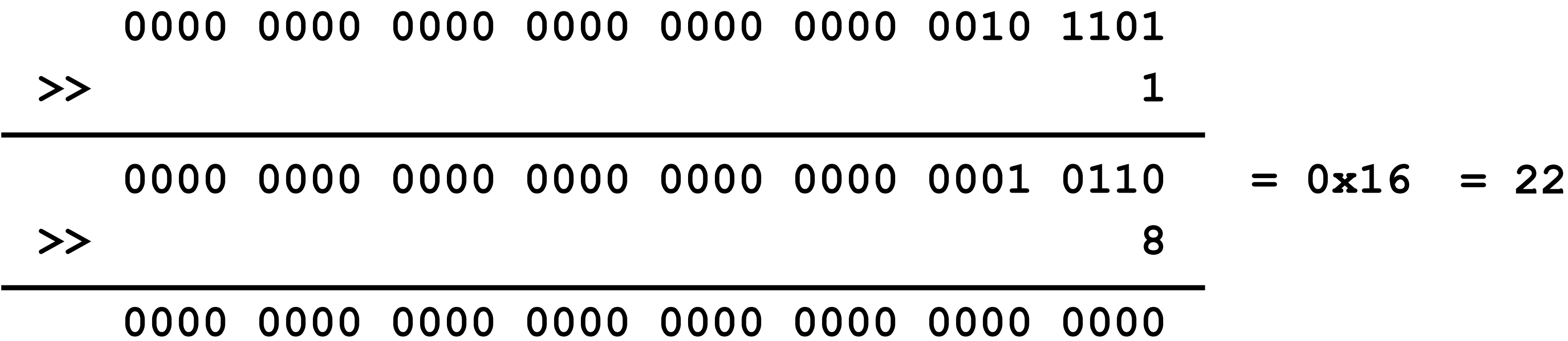
0000 0000 0000 0000 0000 0000 0000 0000

= 0x16 = 22

Unsigned Arithmetics

```
unsigned int n = 45;
```

n	2^n
0	1
1	2
2	4
3	8
4	16
5	32
6	64
7	128



- When shift underflows, the bit is lost.

Bit-packing

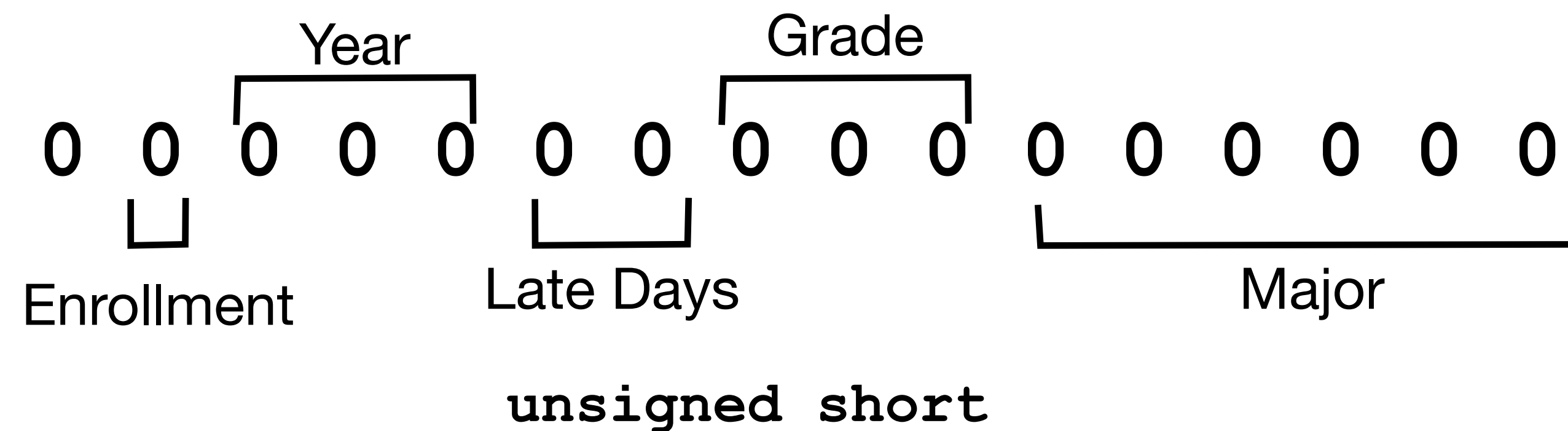
```
struct student {  
    int is_enrolled;    /* is student formally enrolled in this class? */  
    int year;  
    int late_days_used;  
    char letter_grade;  
    char major[32];  
};
```

Bit-packing

- Enrollment: 2 choices (1 bit)
- Year: 5 choices [1, 2, 3, 4, beyond-4] (3 bits)
- Late Days: 4 choices (2 bits)
- Letter grade: 5 choices (3 bits)
- Major: 53 majors (6 bits)
- None of these need 32 bit integers!
- $1 + 3 + 2 + 3 + 6 = 15$, in fact we can fit the entire record inside a short.

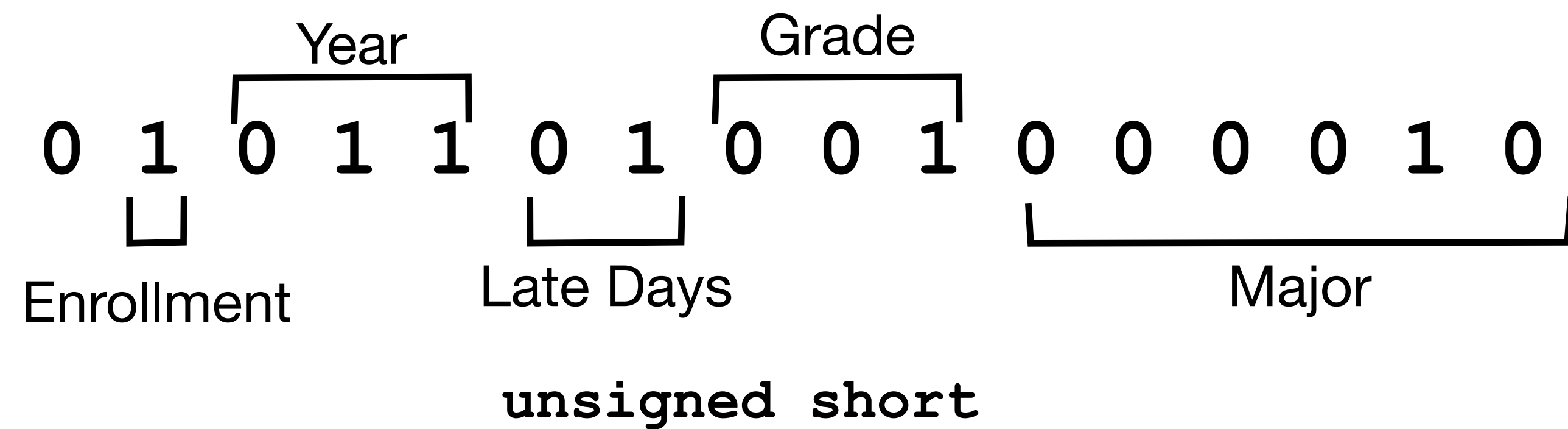
```
struct student {  
    int is_enrolled;  
    int year;  
    int late_days_used;  
    char letter_grade;  
    char major[32];  
};
```

Bit-packing



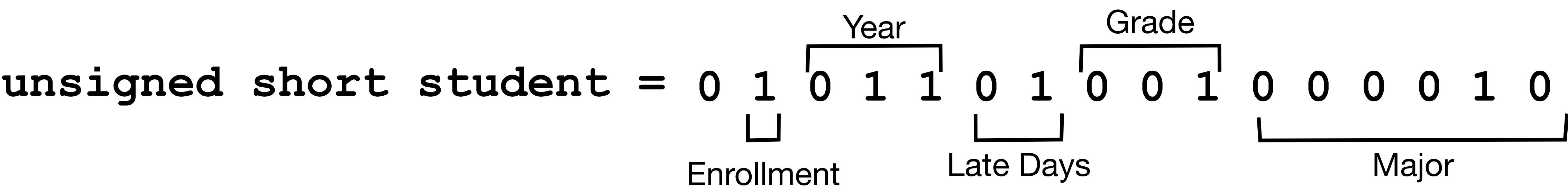
- Enrollment: 2 choices (1 bit)
- Year: 6 choices [0, 1, 2, 3, 4, beyond-4] (3 bits)
- Late Days: 4 choices (2 bits)
- Letter grade: 5 choices (3 bits)
- Major: 53 majors (6 bits)

Bit-packing



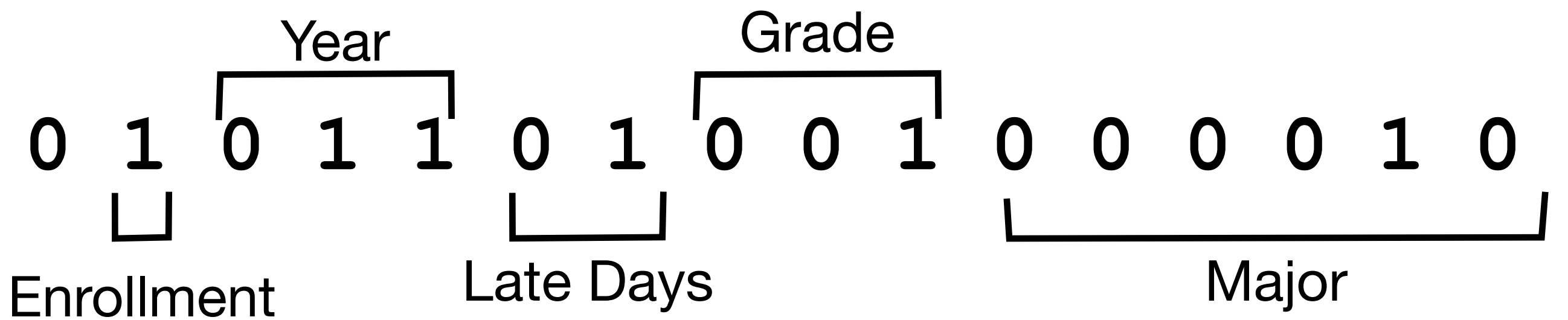
- Enrolled
- Year 3
- 1 late day
- 1st choice of grade
- 2nd choice of major

Bit-packing



Bit-packing

`unsigned short student =` 0 1 0 1 1 0 1 0 0 1 0 0 0 0 1 0



The diagram illustrates the bit fields for the `student` variable. The 16 bits are grouped into five fields: Enrollment (1 bit), Year (3 bits), Late Days (2 bits), Grade (3 bits), and Major (6 bits). The bits are represented as 0s and 1s, with brackets above and below indicating the field boundaries.

Field	Bits
Enrollment	0 1
Year	0 1 1
Late Days	0 1
Grade	0 0 1
Major	0 0 0 0 1 0

`unsigned short mask =`

Bit-packing

unsigned short student = 0 1 0 1 1 0 1 0 0 1 0 0 0 0 1 0

Enrollment Year Late Days Grade Major

The diagram illustrates the bit-packing of student data into a 16-bit unsigned short. The data is divided into five fields: Enrollment (1 bit), Year (3 bits), Late Days (2 bits), Grade (3 bits), and Major (6 bits). The student value is 0101101001000010.

unsigned short mask = 0 0 1 1 1 0 0 0 0 0 0 0 0 0 0 0

The diagram illustrates the bit mask for student data. The mask is 0011100000000000, with the three 1s corresponding to the Year field.

Bit-packing

`unsigned short student =` 0 1 0 1 1 0 1 0 0 1 0 0 0 0 1 0

 ┌─── Year ──┐ ┌─── Grade ──┐

 └──┘ └──┘

 └── Enrollment ─┘ └── Late Days ─┘ └── Major ─┘

`unsigned short mask =` 0 0 1 1 1 0 0 0 0 0 0 0 0 0 0 0

```
int year = mask & student;
```

Bit-packing

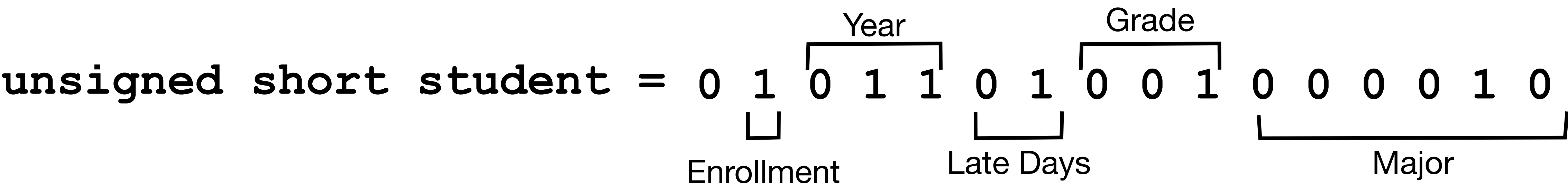
unsigned short student = 0 1 0 1 1 0 1 0 0 1 0 0 0 0 1 0

Enrollment Late Days Year Grade Major

unsigned short mask = 0 0 1 1 1 0 0 0 0 0 0 0 0 0 0 0

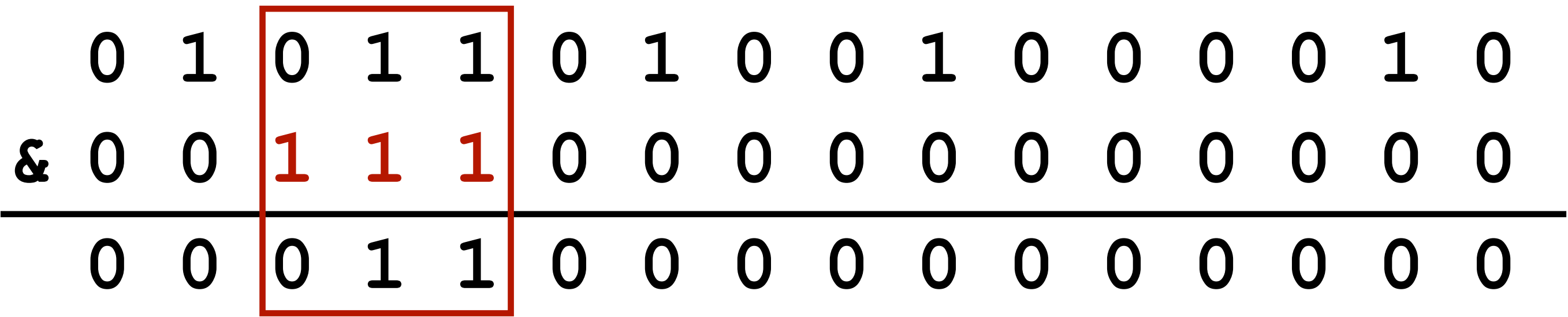
[illegible]

Bit-packing



`unsigned short mask =` 0 0 1 1 1 0 0 0 0 0 0 0 0 0 0 0

```
int year = mask & student;
```



```
year = year >> 11;
```

Bit-packing

unsigned short student = 0 1 0 1 1 0 1 0 0 1 0 0 0 0 1 0

Enrollment Late Days Year Grade Major

unsigned short mask = 0 0 1 1 1 0 0 0 0 0 0 0 0 0 0 0

Diagram illustrating the bit-level operations for the code:

```
int year = mask & student;
```

The diagram shows the bitwise AND operation between the `mask` and `student` variables. The `mask` is represented by the first row of bits (0 1 0 1 1 0 1 0 0 1 0 0 0 0 1 0), and the `student` is represented by the second row of bits (0 0 1 1 1 0 0 0 0 0 0 0 0 0 0 0). The result of the AND operation is shown in the third row (0 0 0 1 1 0 0 0 0 0 0 0 0 0 0 0). The bits 1 1 1 in the student row and 0 1 1 in the result row are highlighted with red boxes.

```
year = year >> 11;
```

The diagram shows the right shift operation by 11 bits. The original result (0 0 0 1 1 0 0 0 0 0 0 0 0 0 0 0) is shifted to the right by 11 positions, resulting in the final value (0 0 0 0 0 0 0 0 0 0 0 0 0 1 1 0). The original result and the final value are shown in the fourth and fifth rows, respectively. The final value's last three bits (1 1 0) are highlighted with a red box.

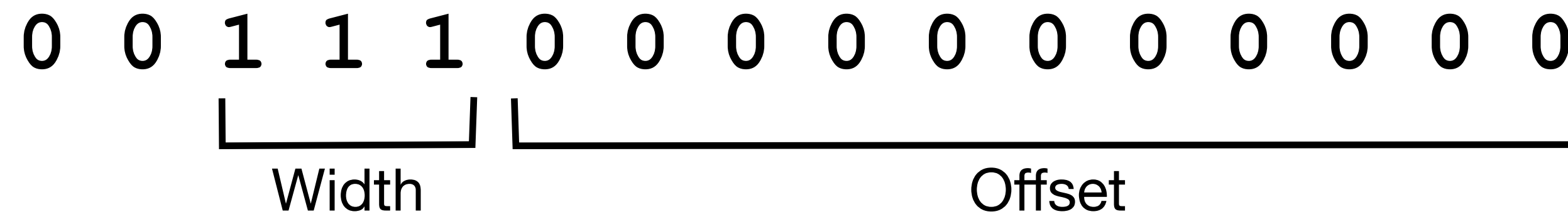
Bit-packing

0 0 1 1 1 0 0 0 0 0 0 0 0 0 0 0

```
unsigned short mask = 14336;
```

```
unsigned short mask = 0x7 << 11;
```

Bit-packing



```
unsigned short mask = 14336;
```

```
unsigned short mask = 0x7 << 11;
```

```
unsigned short mask = ~0u >> (16 - width) << offset;
```

Future Note: we need to write 0u instead of 0 because by default, 0 is signed, so >> is going to perform a signed right shift.

Bit-packing

0 0 1 1 1 0 0 0 0 0 0 0 0 0 0 0

└──┬──────────────────────────┘

Width Offset

```
unsigned short mask = ~0 >> (16 - width) << offset;
```

~ 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0

Bit-packing

0 0 1 1 1 0 0 0 0 0 0 0 0 0 0 0

└───┬──────────────────────────┘

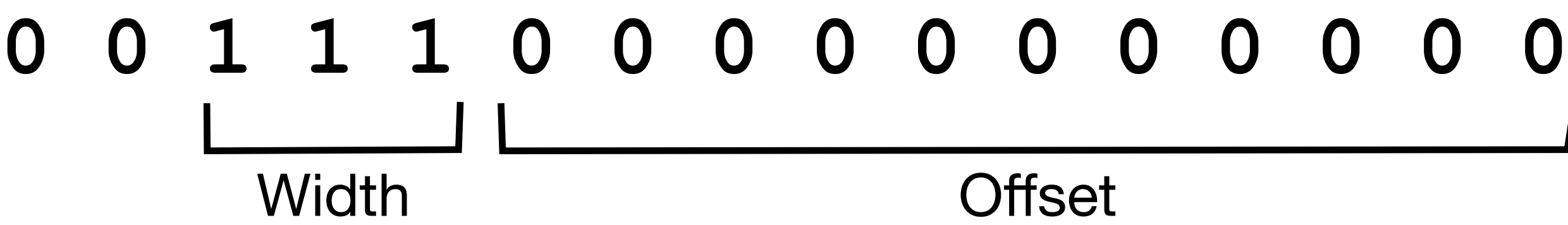
Width Offset

```
unsigned short mask = ~0 >> (16 - width) << offset;
```

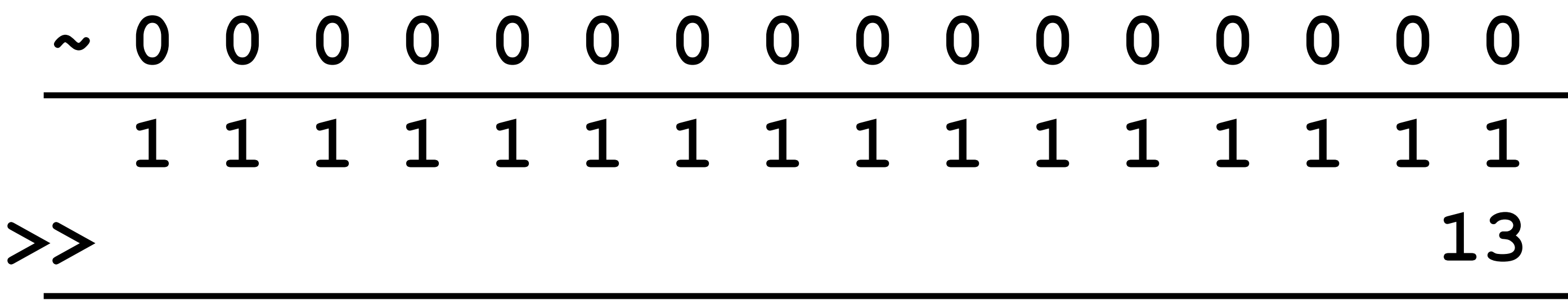
~ 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0

1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1

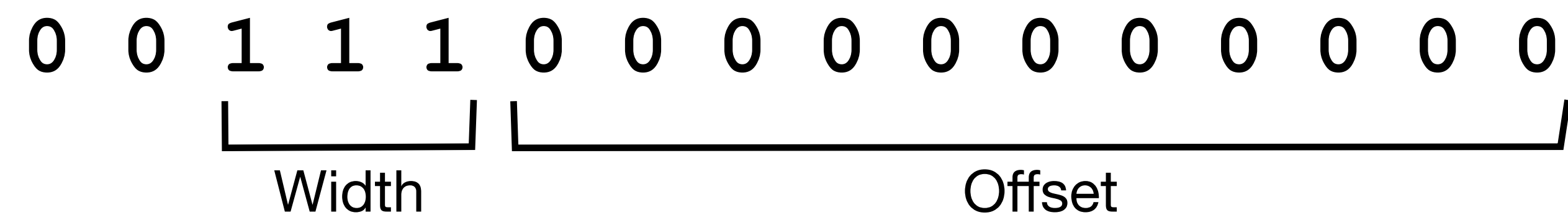
Bit-packing



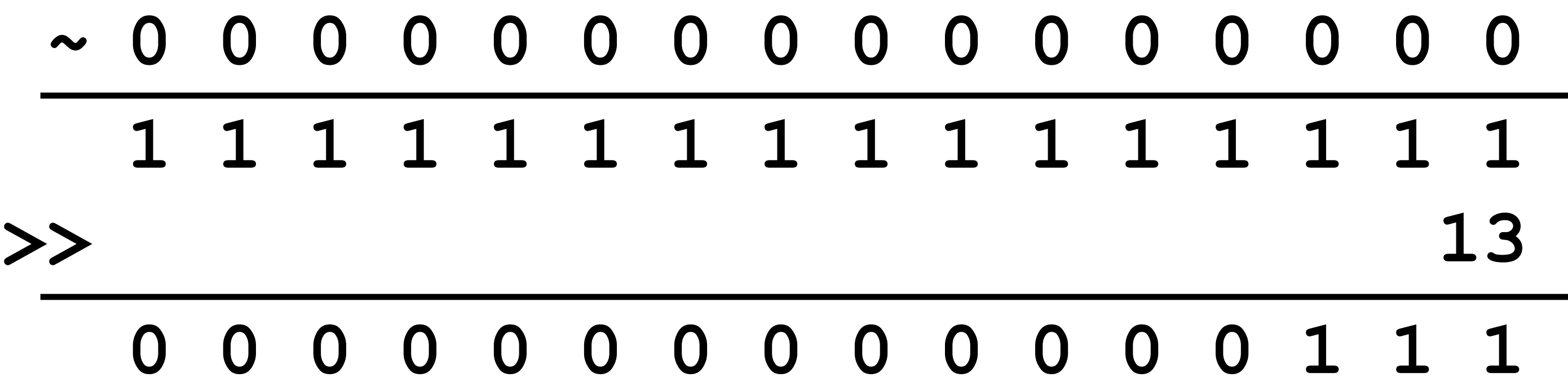
```
unsigned short mask = ~0 >> (16 - width) << offset;
```



Bit-packing



```
unsigned short mask = ~0 >> (16 - width) << offset;
```



Bit-packing

0 0 1 1 1 0 0 0 0 0 0 0 0 0 0 0

└──┬──────────────────────────┘

Width Offset

```
unsigned short mask = ~0 >> (16 - width) << offset;
```

~ 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0

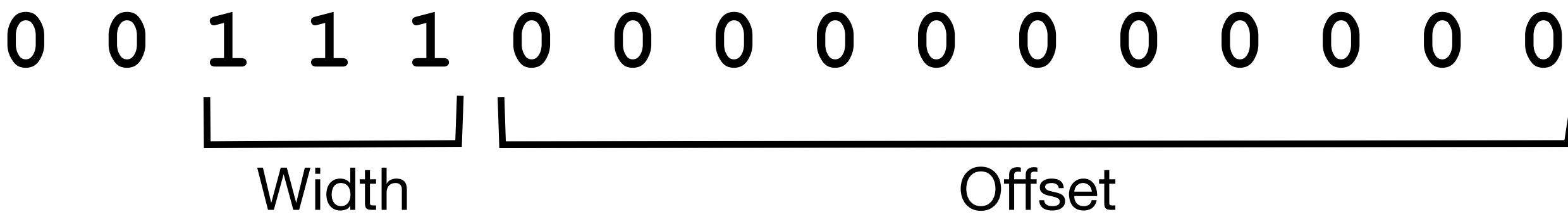
1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1

>> 13

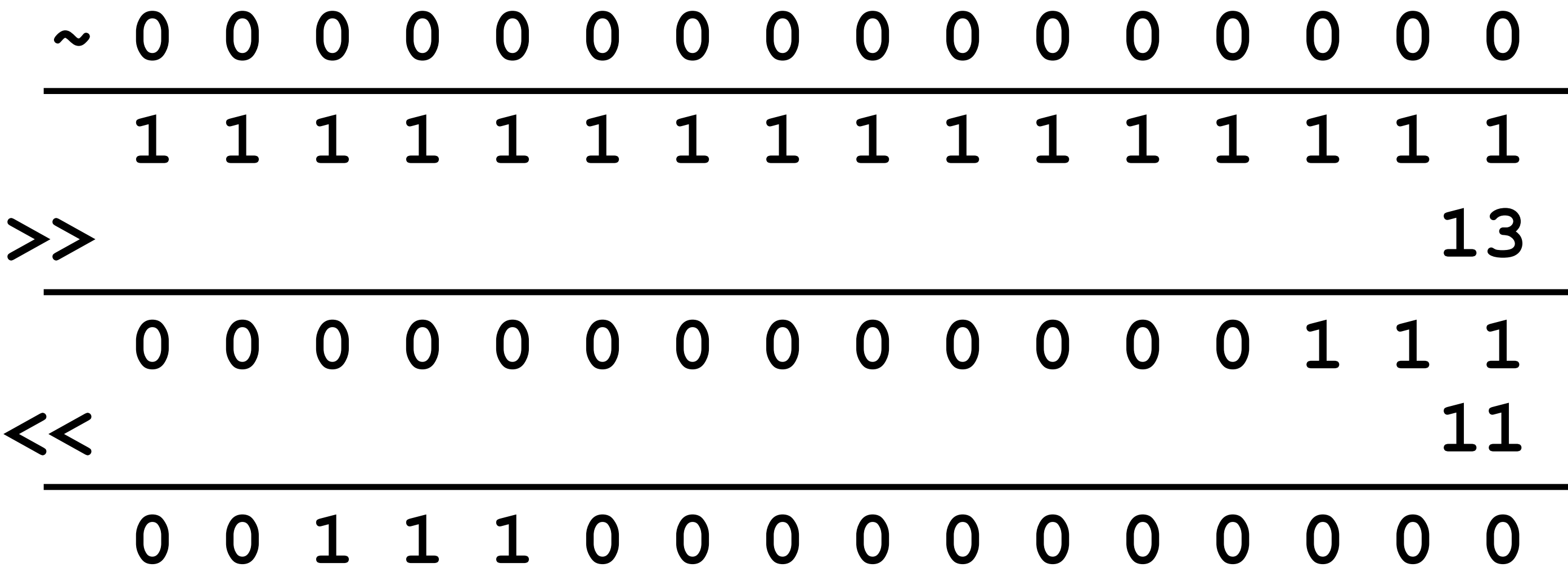
0 0 0 0 0 0 0 0 0 0 0 0 0 1 1 1

<< 11

Bit-packing



```
unsigned short mask = ~0 >> (16 - width) << offset;
```



Bit-packing

```
unsigned short student = ...;  
unsigned short year_mask = ~0 >> (16 - 3) << 11;  
  
unsigned int year = (student & year_mask) >> 11;
```

Bit-packing

`unsigned short student =` 0 1 0 1 1 0 1 0 0 1 0 0 0 0 1 0

Enrollment Year Late Days Grade Major

`unsigned short mask =` 0 0 1 1 1 0 0 0 0 0 0 0 0 0 0 0

`unsigned short year =` 0 0 0 0 0 0 0 0 0 0 0 0 0 0 1 0 0

Bit-packing

`unsigned short student =` 0 1 0 1 1 0 1 0 0 1 0 0 0 0 1 0

Enrollment Year Late Days Grade Major

`unsigned short mask =` 0 0 1 1 1 0 0 0 0 0 0 0 0 0 0 0

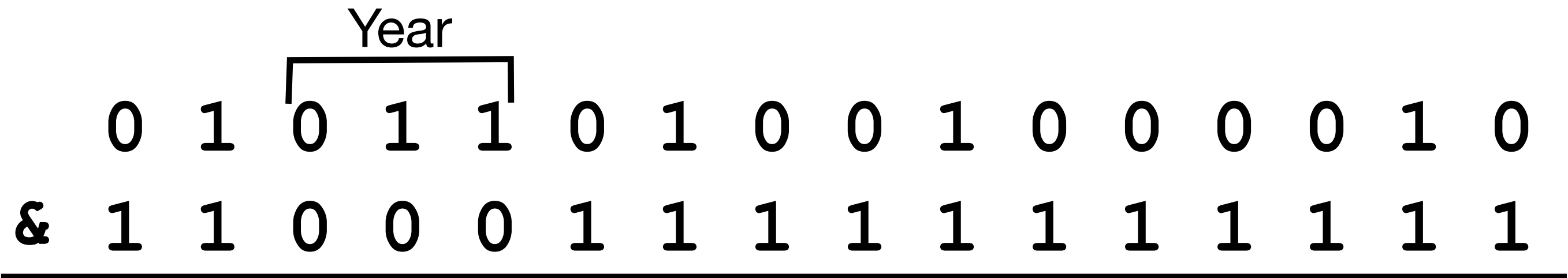
`unsigned short year =` 0 0 0 0 0 0 0 0 0 0 0 0 0 0 1 0 0

`year = year << 11;` 0 0 1 0 0 0 0 0 0 0 0 0 0 0 0 0

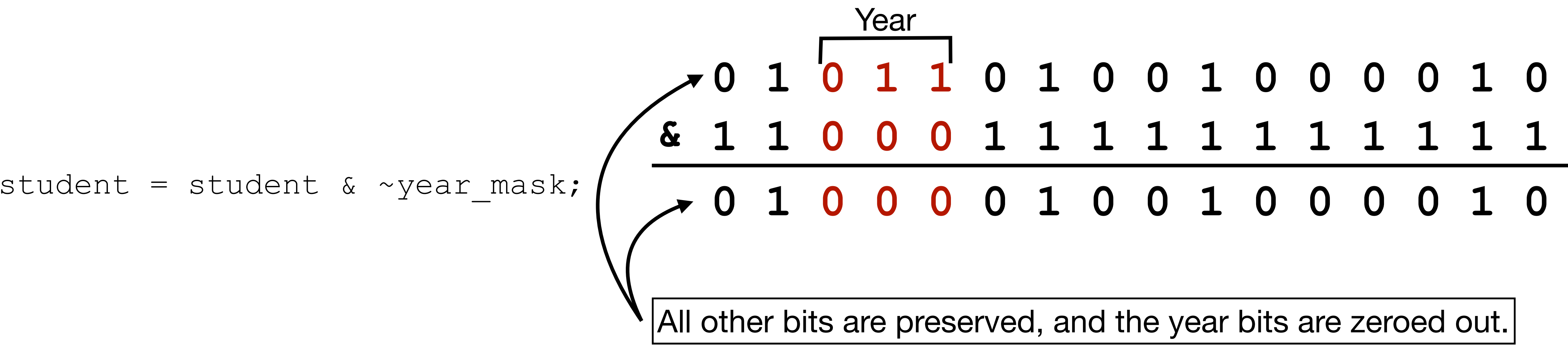
??

Bit-packing

```
student = student & ~year_mask;
```



Bit-packing



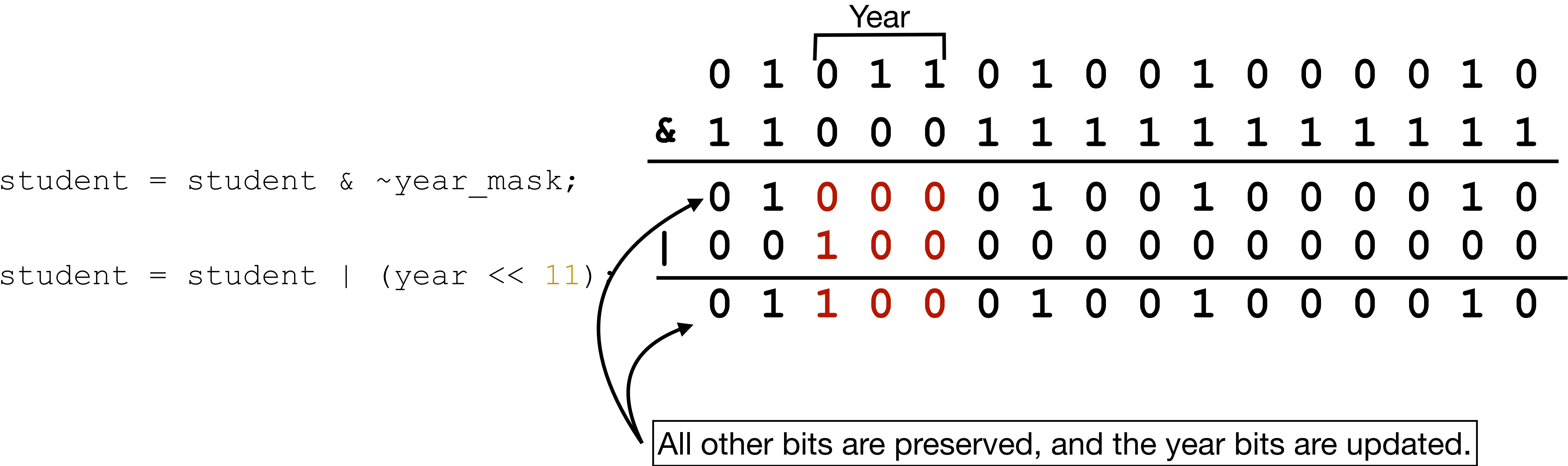
Bit-packing

```
student = student & ~year_mask;
```

```
student = student | (year << 11);
```

				Year												
	0	1	0	1	1	0	1	0	0	1	0	0	0	0	1	0
&	1	1	0	0	0	1	1	1	1	1	1	1	1	1	1	1
	0	1	0	0	0	0	1	0	0	1	0	0	0	0	1	0
	0	0	1	0	0	0	0	0	0	0	0	0	0	0	0	0

Bit-packing



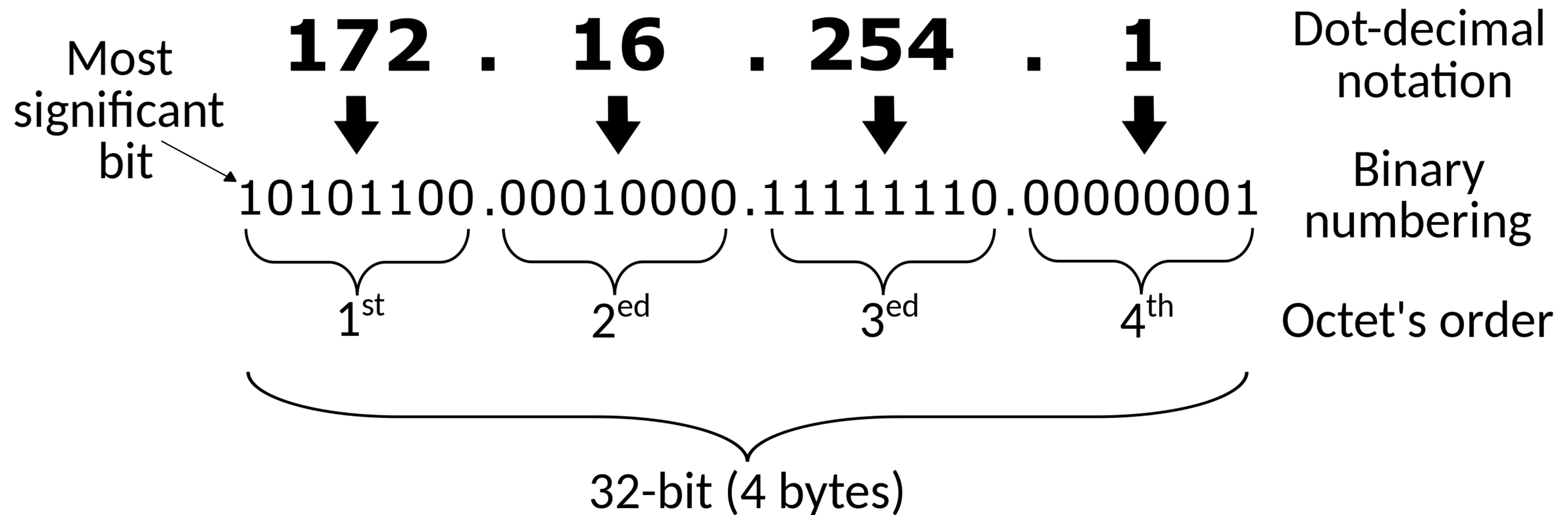
Bit-packing

Summary

- A mask has the bit fields of interest set 1 and all others 0.
- `& mask` gets the bits
- `& ~mask` clears out the bits
- `~0u >> (16 - width) << offset` creates a 16-bit mask with `width` and `offset`.
- After the bits are zeroed out, `|` is used to write new bits in the field.
- Does anyone use this irl?

Bit-packing

Example: IP Address



Bit-packing

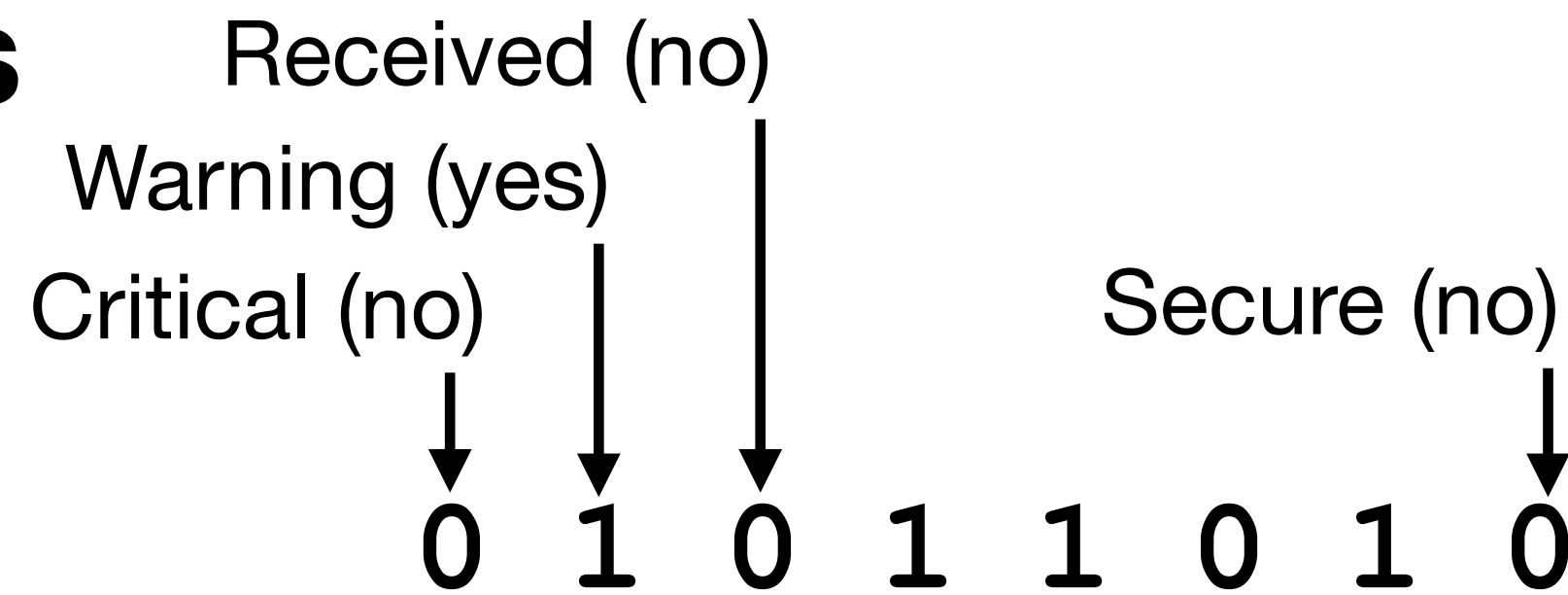
Example: IP Address

IP address	192.168.1.2
Subnet mask	255.255.255.0
Router	192.168.1.1

What is this mask trying to get?

Bit-packing

Example: Flag Bits

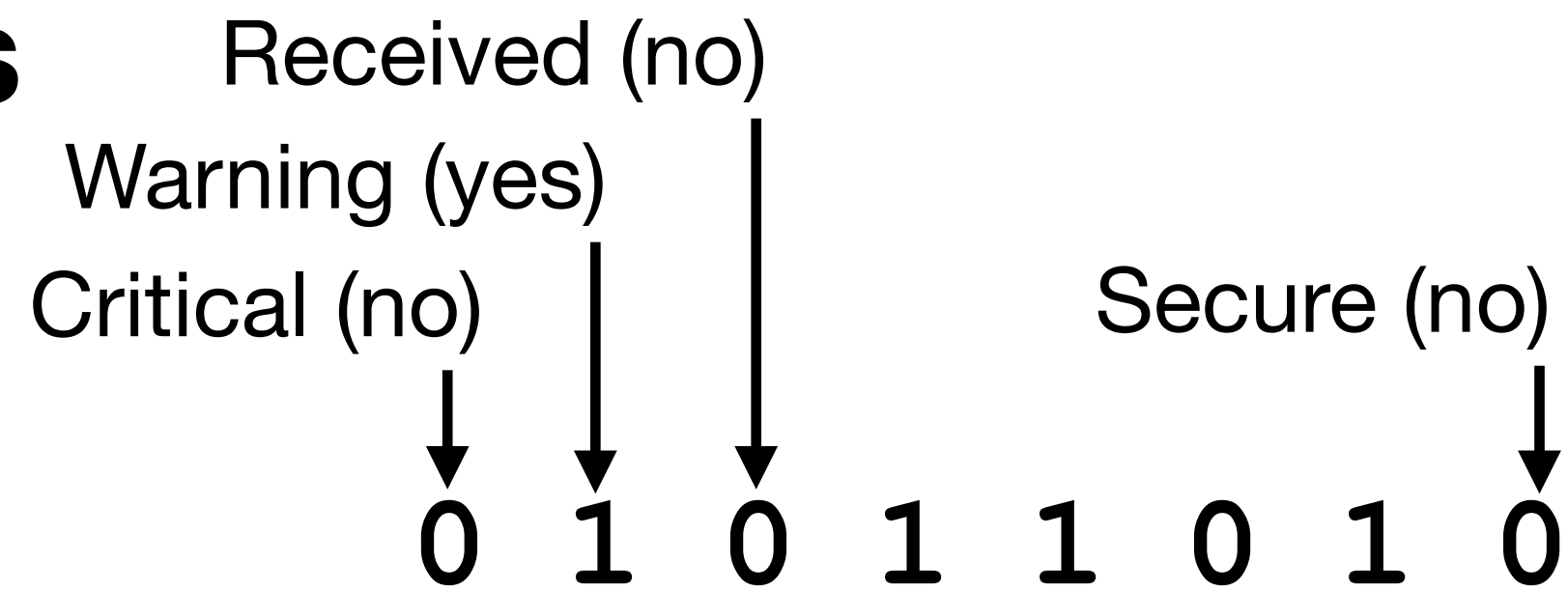


Each bit answers one yes/no question -- 8 Booleans in a byte!

```
const uint8_t CRITICAL = 0x80;    ->100000000
const uint8_t WARNING  = 0x40;    ->010000000
const uint8_t RECEIVED = 0x20;    ->001000000
...
const uint8_t SECURE   = 0x01;    ->000000001
```

Bit-packing

Example: Flag Bits



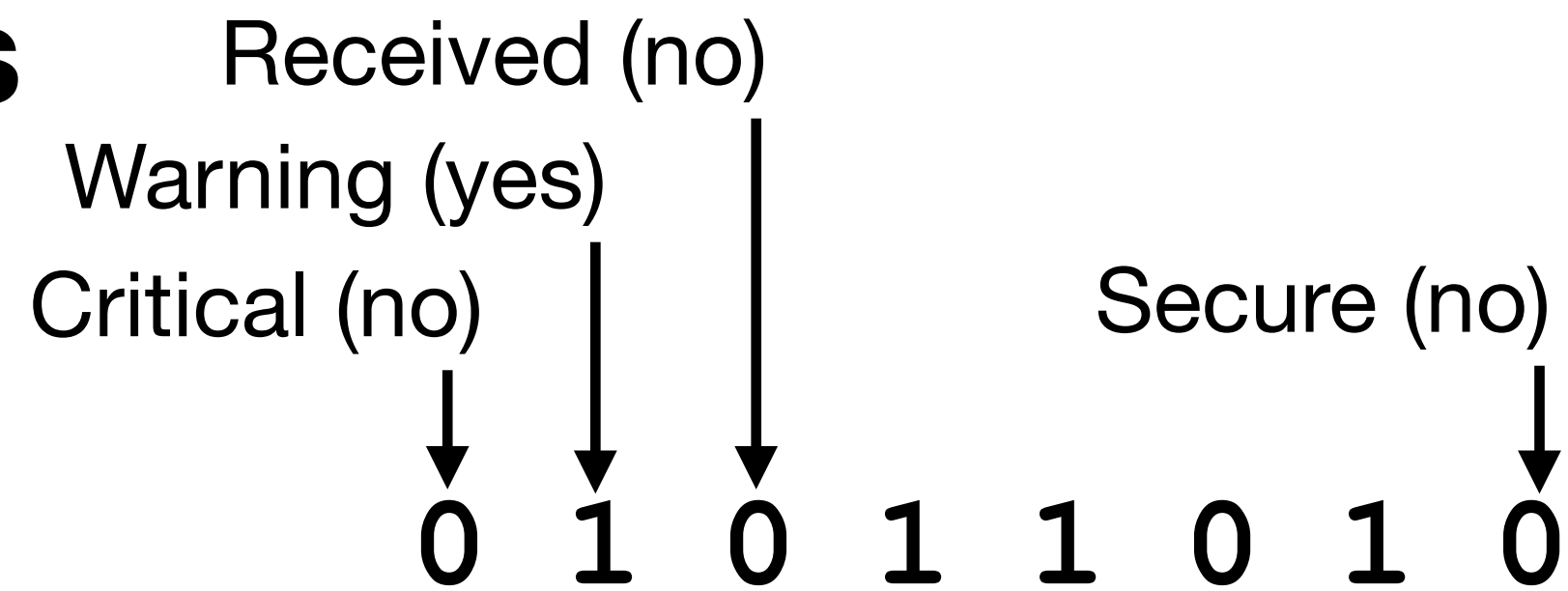
```
const uint8_t CRITICAL = 0x80;  -> 10000000
const uint8_t WARNING  = 0x40;  -> 01000000
const uint8_t RECEIVED = 0x20;  -> 00100000
...
const uint8_t SECURE   = 0x01;  -> 00000001
```

```
if (flags & CRITICAL) { ... }
```

This is non-zero if and only if the first bit is set.

Bit-packing

Example: Flag Bits



```
const uint8_t CRITICAL = 0x80;    -> 100000000
const uint8_t WARNING  = 0x40;    -> 010000000
const uint8_t RECEIVED = 0x20;    -> 001000000
...
const uint8_t SECURE   = 0x01;    -> 000000001
```

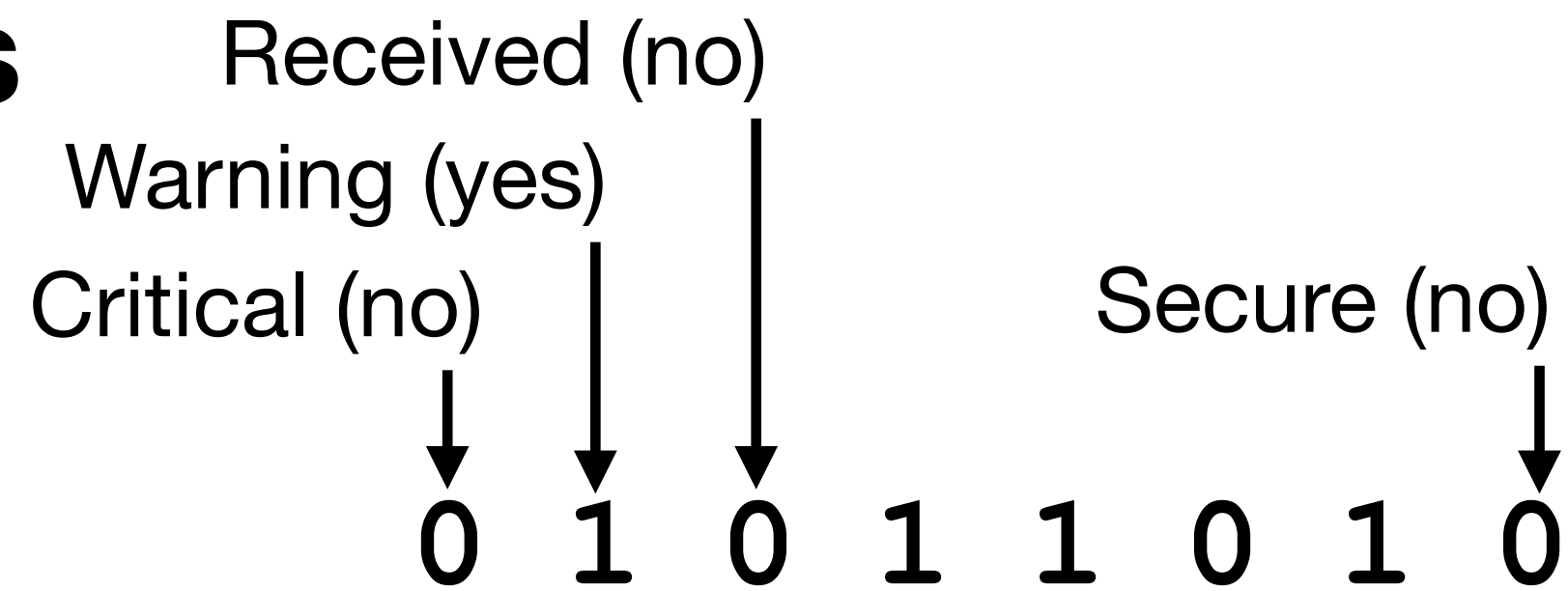
```
if (flags & (CRITICAL | SECURE)) { ... }
```

10000001

This is non-zero if either the first bit or the last bit is set.

Bit-packing

Example: Flag Bits



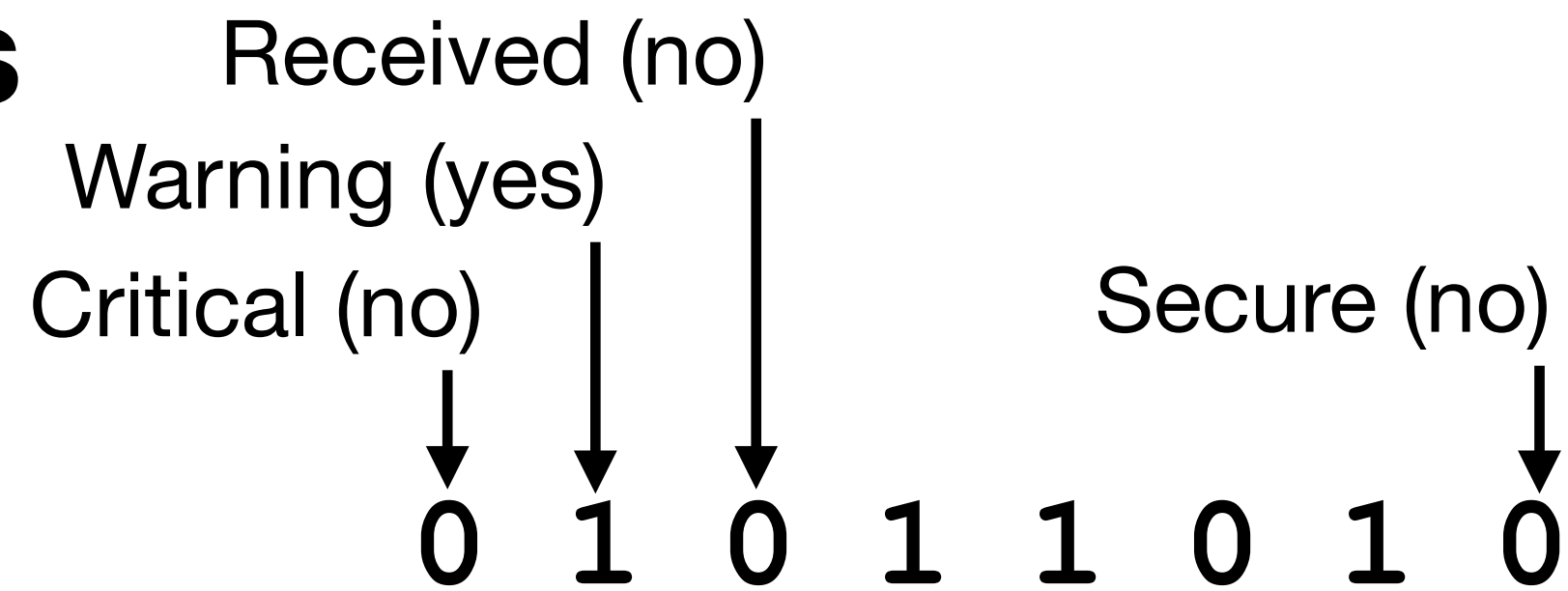
```
const uint8_t CRITICAL = 0x80;  -> 100000000
const uint8_t WARNING  = 0x40;  -> 010000000
const uint8_t RECEIVED = 0x20;  -> 001000000
...
const uint8_t SECURE   = 0x01;  -> 000000001
```

```
flags |= RECEIVED;
```

This sets the RECEIVED bit.

Bit-packing

Example: Flag Bits



```
const uint8_t CRITICAL = 0x80;  -> 100000000
const uint8_t WARNING  = 0x40;  -> 010000000
const uint8_t RECEIVED = 0x20;  -> 001000000
...
const uint8_t SECURE   = 0x01;  -> 000000001
```

```
flags &= ~SECURE;
```

This unsets the SECURE bit.

Interlude: Command-Line Arguments

and ++, --

Negative Numbers

What is this?

1110 1001

- 233
- Ascii é
- Yes Yes Yes No Yes No No Yes
- Sound wave sample with some magnitude
- A pixel intensity of 91.31%
- ...
- It's one of the 256 choices, whatever that is.

Negative Numbers

- If we have 32 bits, we have 2^{32} choices
 - unsigned: 2^{32} choices of non-negative numbers -- 0 through $2^{32} - 1$
 - signed int: 2^{32} choices of both negative and positive numbers *about* -2^{31} through 2^{31}

Negative Numbers

Attempt 1

- The first bit is the sign, 0 -> positive, 1 -> negative

5 = 0000 0000 0000 0000 0000 0000 0000 0011

-5 = 1000 0000 0000 0000 0000 0000 0000 0011

- Range?
 - $[-(2^{31} - 1), 2^{31} - 1]$

Negative Numbers

Attempt 1

- The first bit is the sign, 0 -> positive, 1 -> negative

5 = 0000 0000 0000 0000 0000 0000 0000 0011

-5 = 1000 0000 0000 0000 0000 0000 0000 0011

- What's wrong with this?

0000 0000 0000 0000 0000 0000 0000 0000

1000 0000 0000 0000 0000 0000 0000 0000

Negative Numbers

Attempt 1

- The first bit is the sign, 0 -> positive, 1 -> negative

5 = 0000 0000 0000 0000 0000 0000 0000 0011

-5 = 1000 0000 0000 0000 0000 0000 0000 0011

- What's wrong with this?

0000 0000 0000 0000 0000 0000 0000 0000

1000 0000 0000 0000 0000 0000 0000 0000

- We have +0 and -0 !?
 - Waste a combination, equality check is complicated +0 == -0

Negative Numbers

Attempt 1

- The first bit is the sign, 0 -> positive, 1 -> negative

5 = 0000 0000 0000 0000 0000 0000 0000 0011

-5 = 1000 0000 0000 0000 0000 0000 0000 0011

- What's wrong with this?

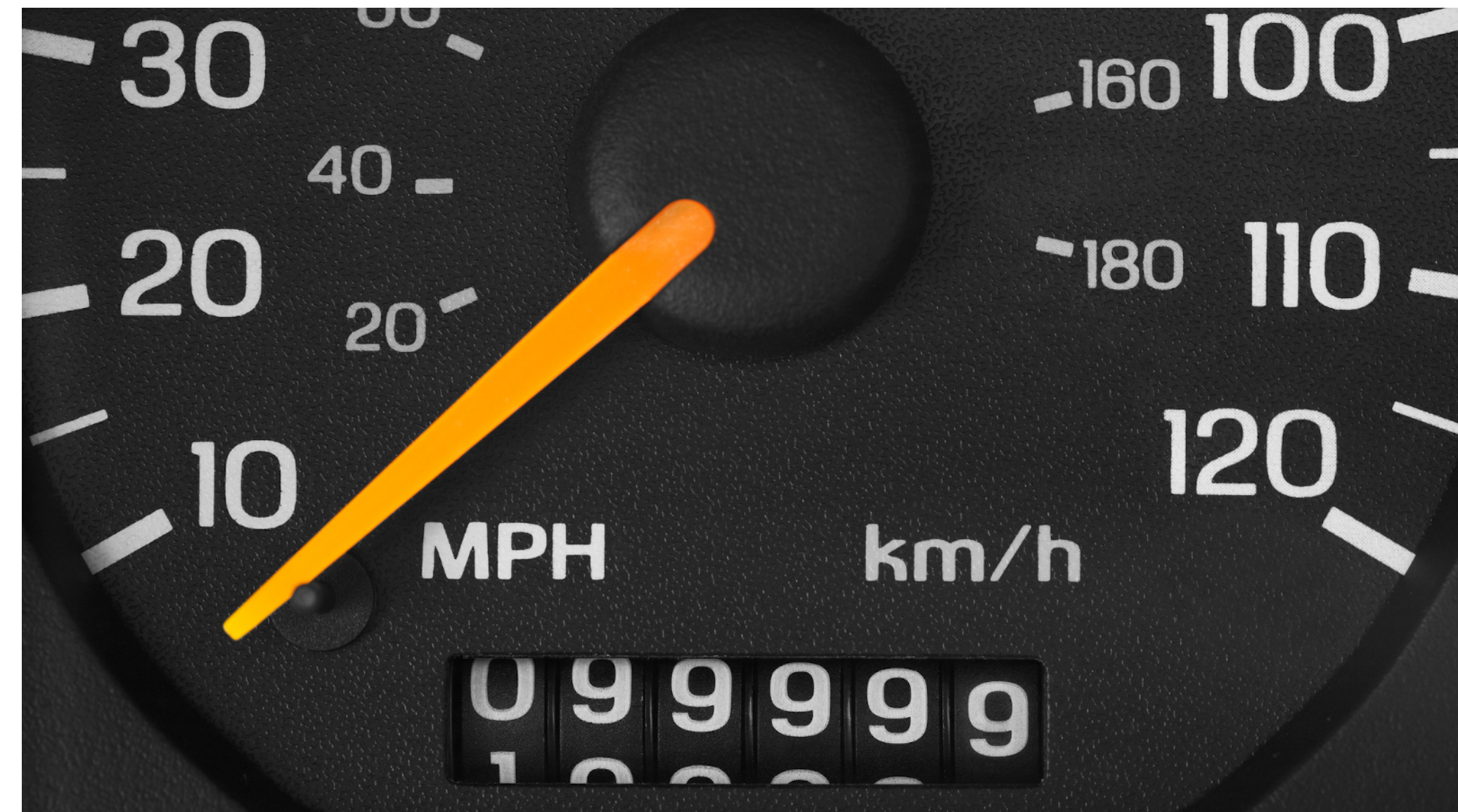
0000 0000 0000 0000 0000 0000 0000 0011

+ 1000 0000 0000 0000 0000 0000 0000 0011

- 5 + (-5) is not 0 :(((
 - Need special circuit for adding signed numbers and unsigned numbers

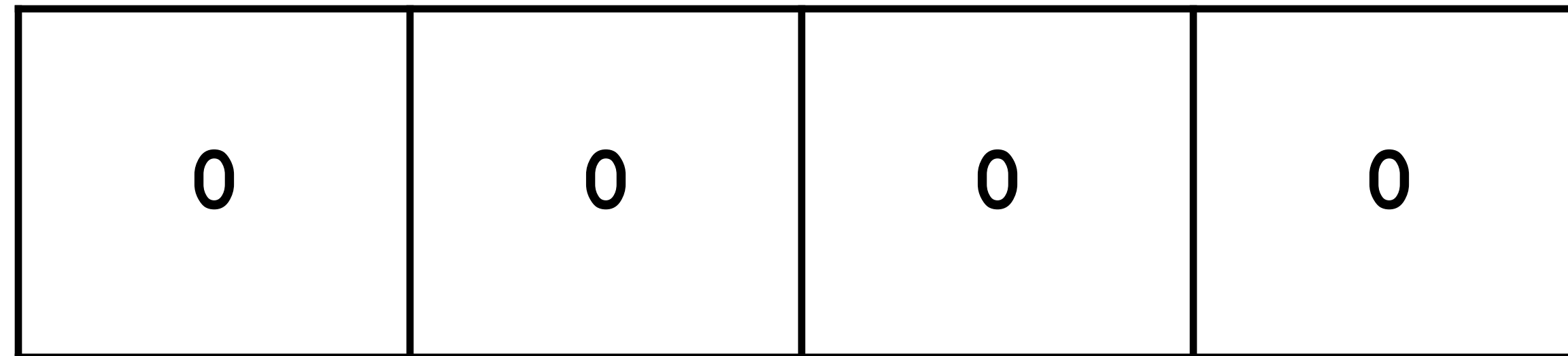
Negative Numbers

Better Idea: Two's Complement



Negative Numbers

Better Idea: Two's Complement



Negative Numbers

Better Idea: Two's Complement

0	0	0	1
---	---	---	---

Negative Numbers

Better Idea: Two's Complement

0	0	0	2
---	---	---	---

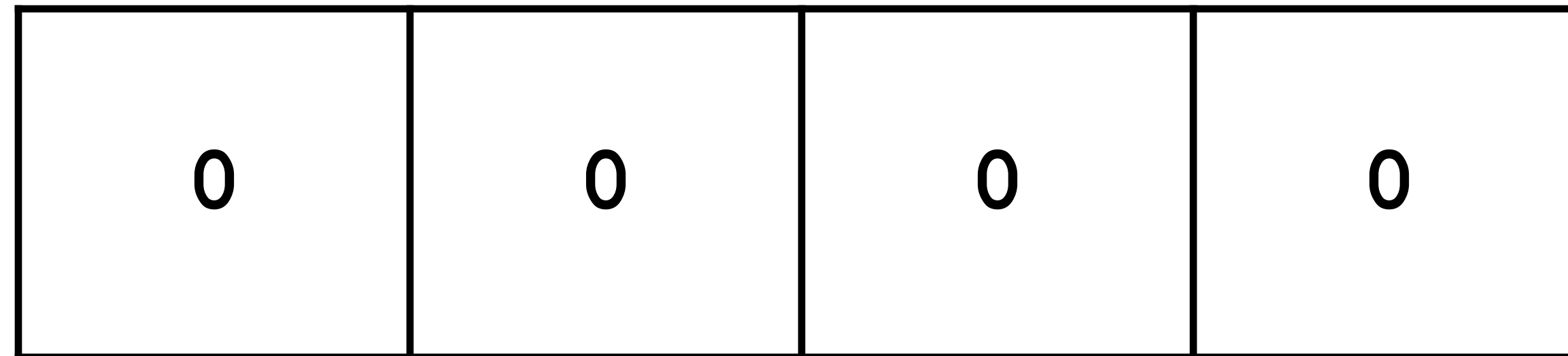
Negative Numbers

Better Idea: Two's Complement

0	0	0	1
---	---	---	---

Negative Numbers

Better Idea: Two's Complement



Negative Numbers

Better Idea: Two's Complement

9	9	9	9
---	---	---	---

Negative Numbers

Better Idea: Two's Complement

9	9	9	8
---	---	---	---

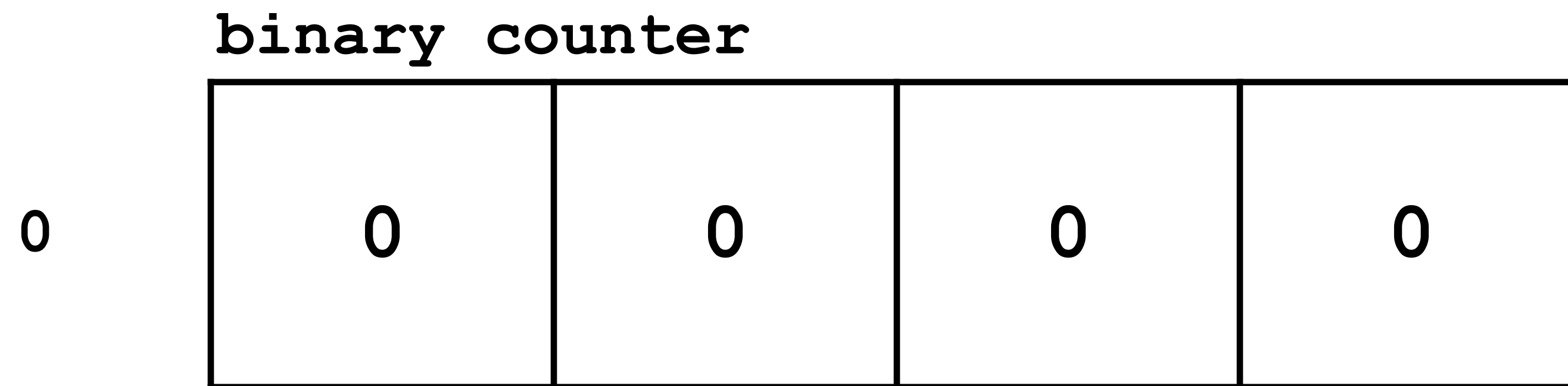
Negative Numbers

Better Idea: Two's Complement

9	9	9	7
---	---	---	---

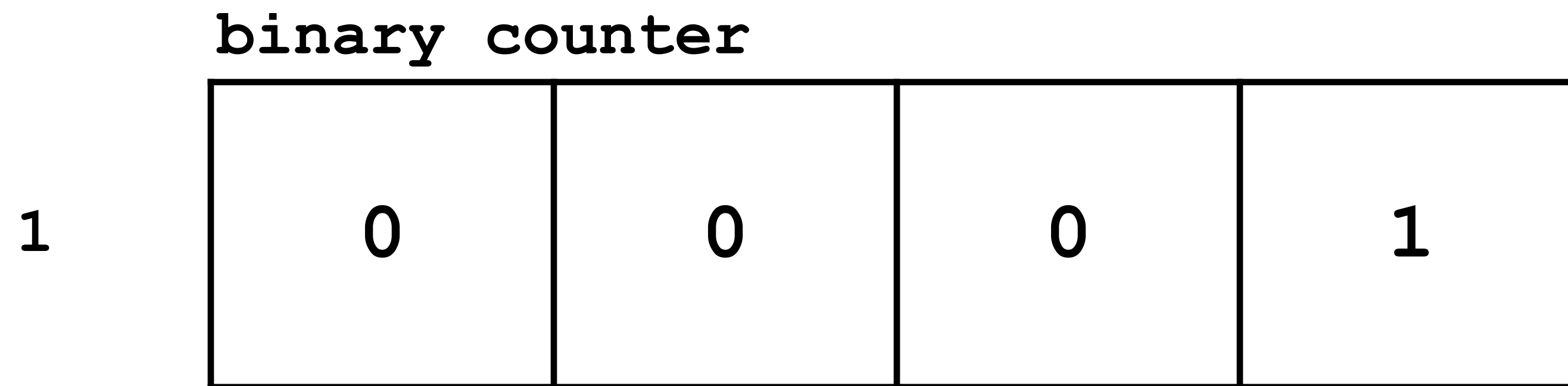
Negative Numbers

Better Idea: Two's Complement



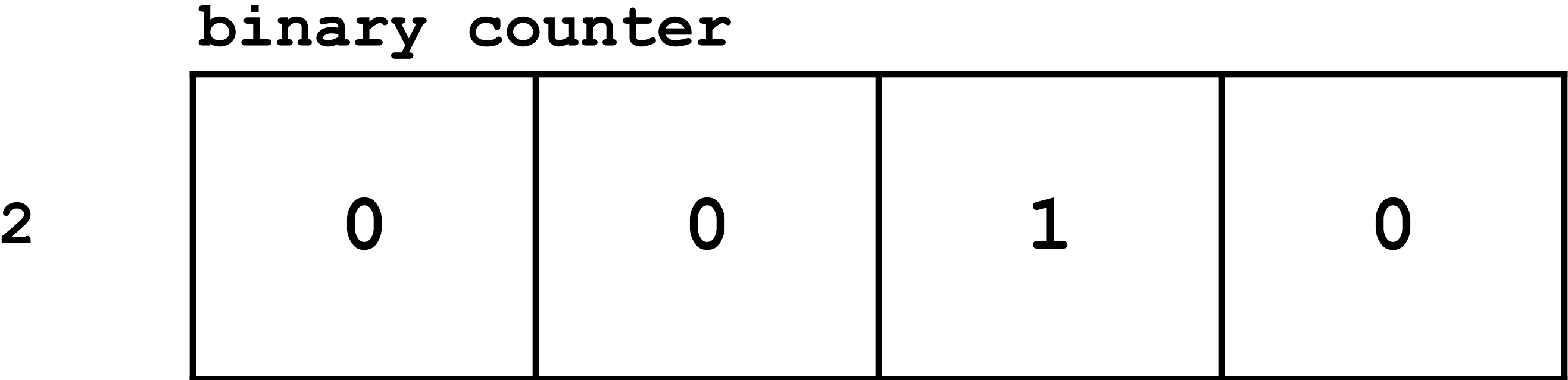
Negative Numbers

Better Idea: Two's Complement



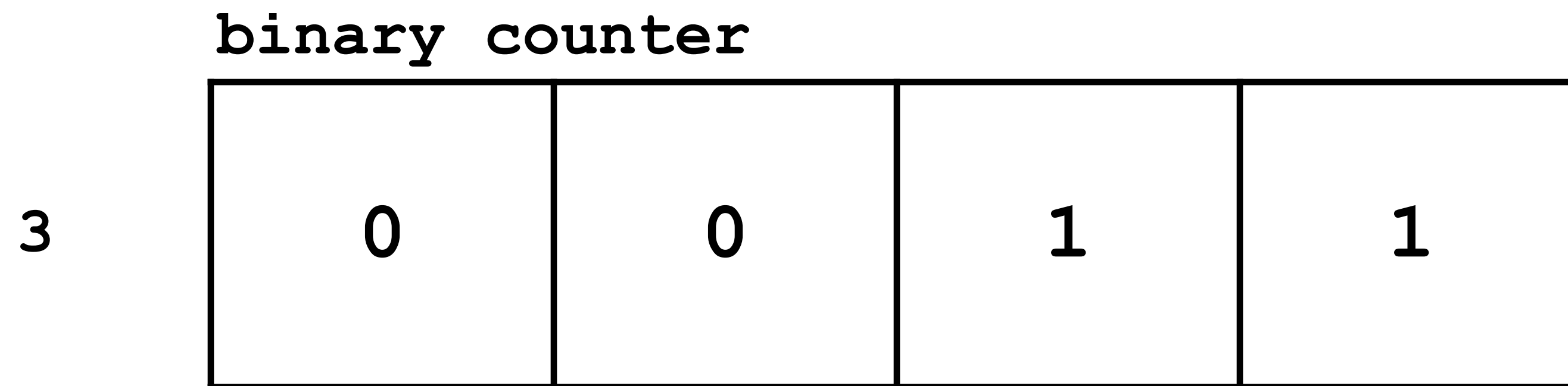
Negative Numbers

Better Idea: Two's Complement



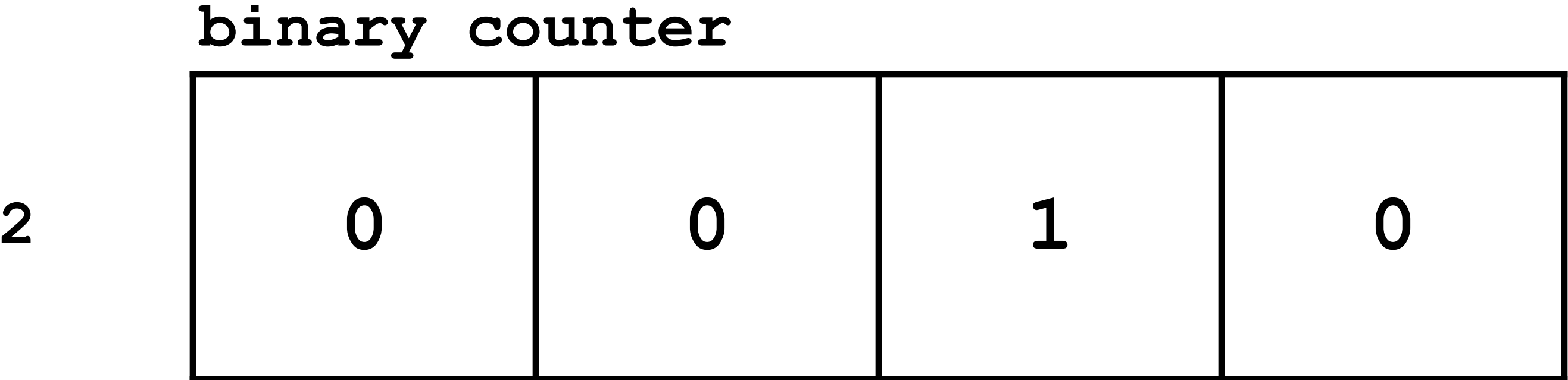
Negative Numbers

Better Idea: Two's Complement



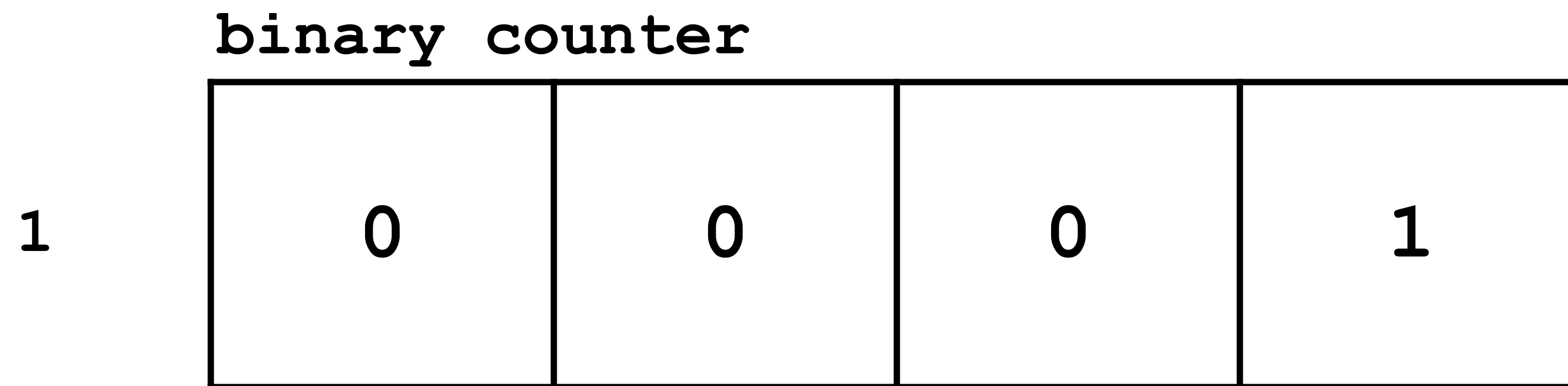
Negative Numbers

Better Idea: Two's Complement



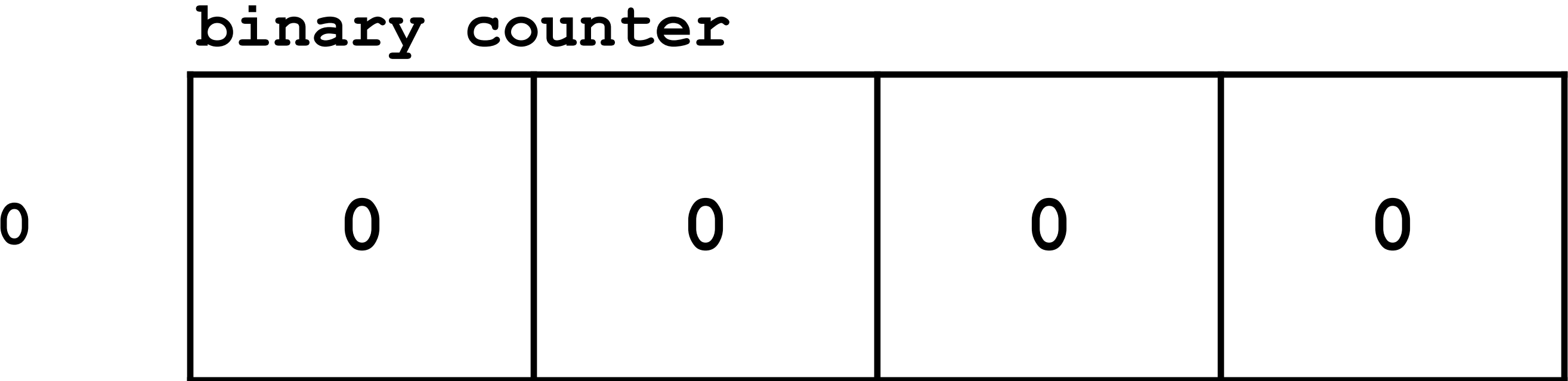
Negative Numbers

Better Idea: Two's Complement



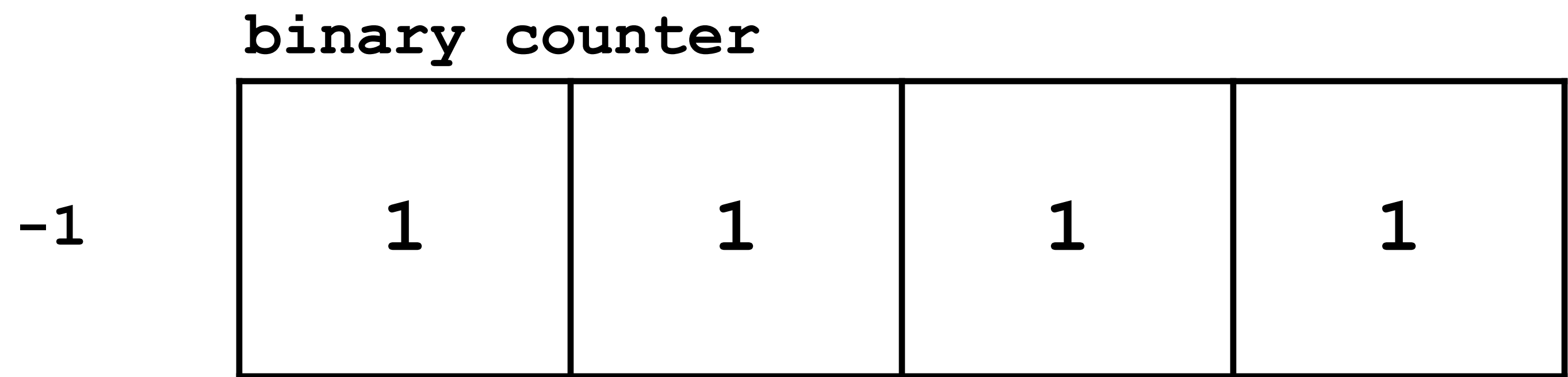
Negative Numbers

Better Idea: Two's Complement



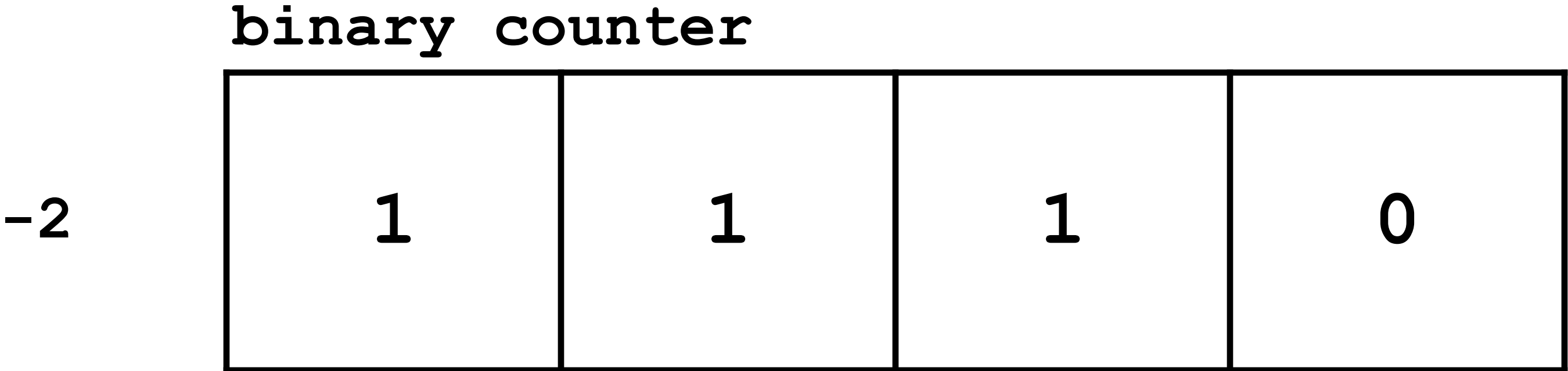
Negative Numbers

Better Idea: Two's Complement



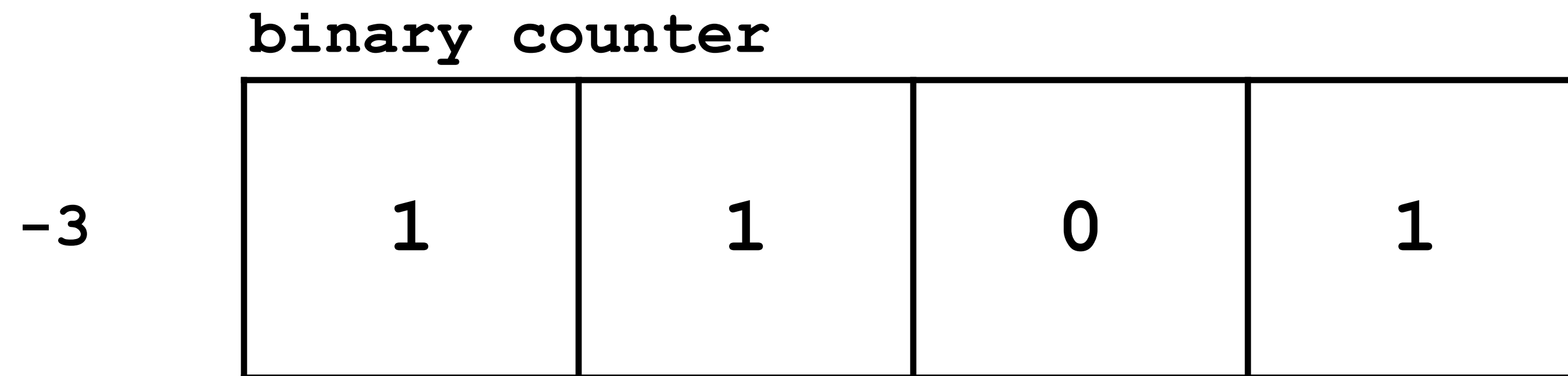
Negative Numbers

Better Idea: Two's Complement



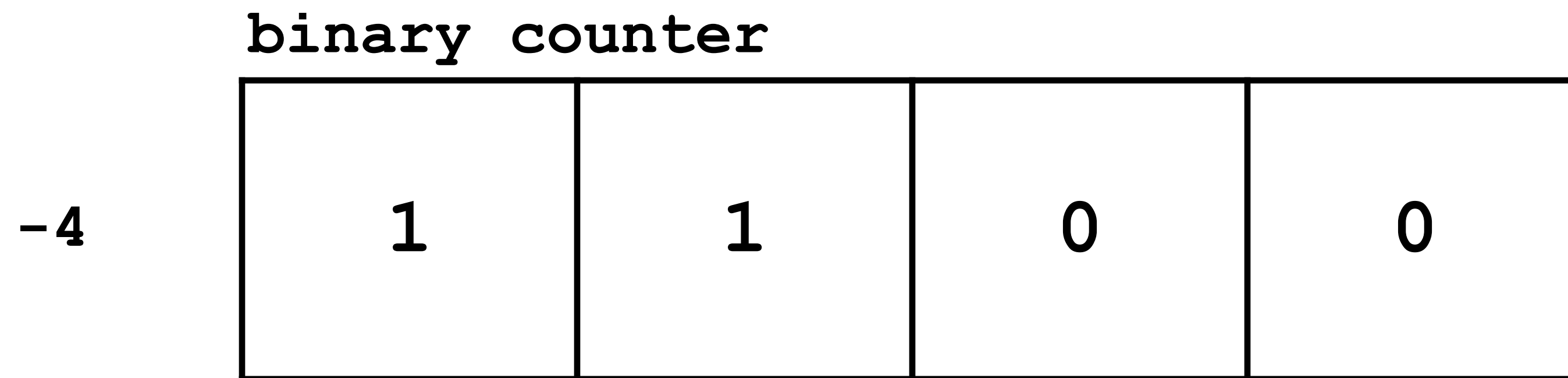
Negative Numbers

Better Idea: Two's Complement



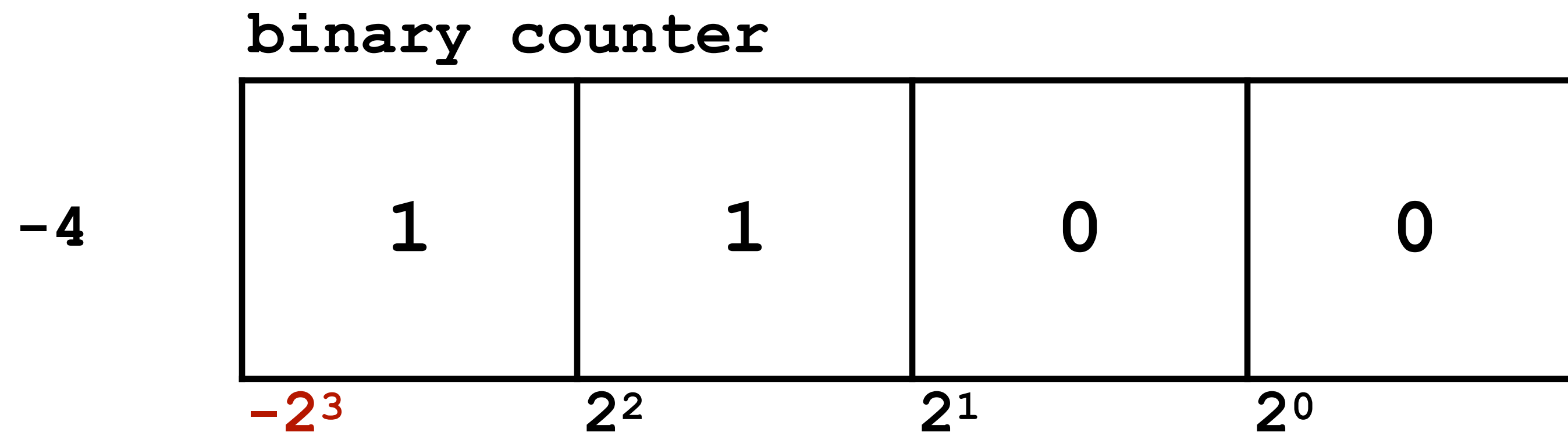
Negative Numbers

Better Idea: Two's Complement



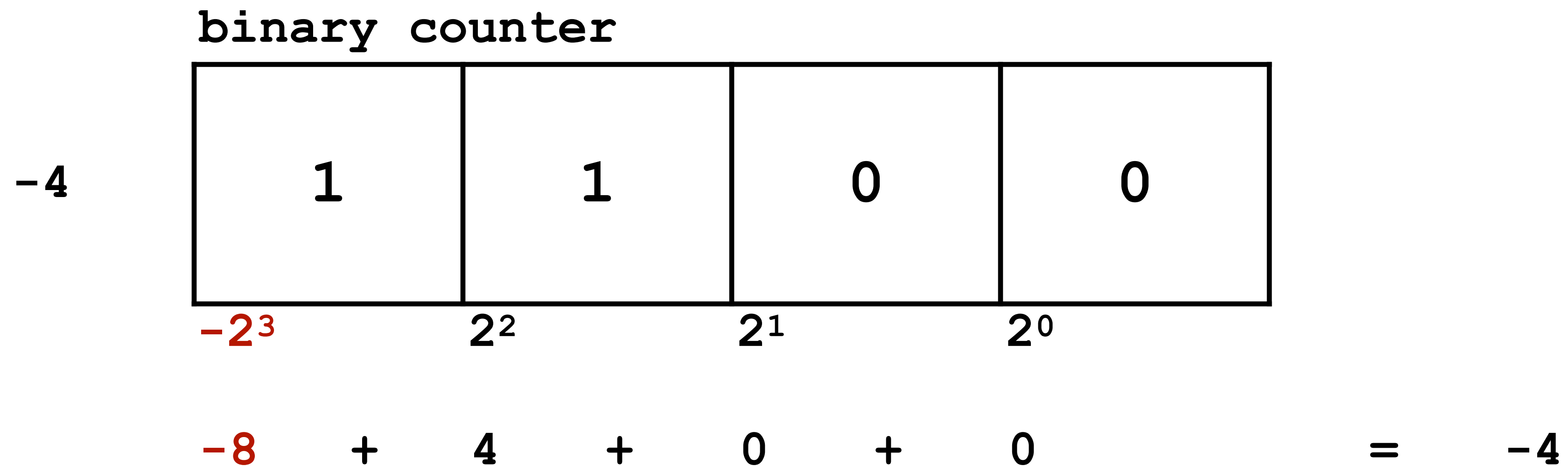
Negative Numbers

Two's complement: two ways



Negative Numbers

Two's complement: two ways



- The most significant bit is negative.

Negative Numbers

Two's complement: two ways

n	2^n
0	1
1	2
2	4
3	8
4	16
5	32
6	64
7	128

1110 1001

- Unsigned: 233
- What does this represent if it is a signed integer?
- **-128** + 64 + 32 + 8 + 1 = -23

Negative Numbers

Two's complement: two ways

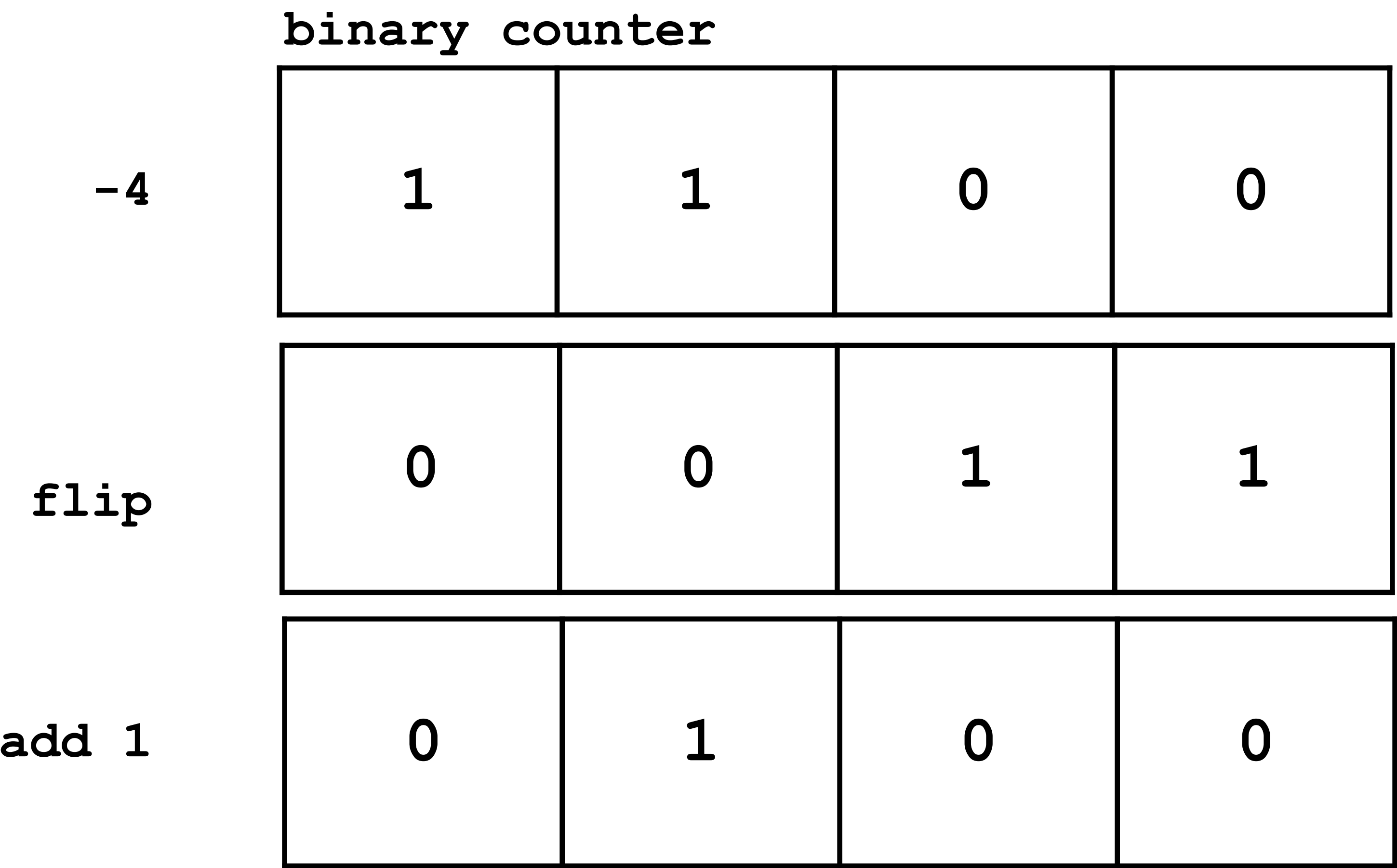
n	2^n
0	1
1	2
2	4
3	8
4	16
5	32
6	64
7	128

1110 1001

- How do hardwares negate the sign?
 - Flip all the bits and add one
- Flip all the bits: 0001 0110
- add one: 0001 0111
- : 16 + 4 + 2 + 1 = 23

Negative Numbers

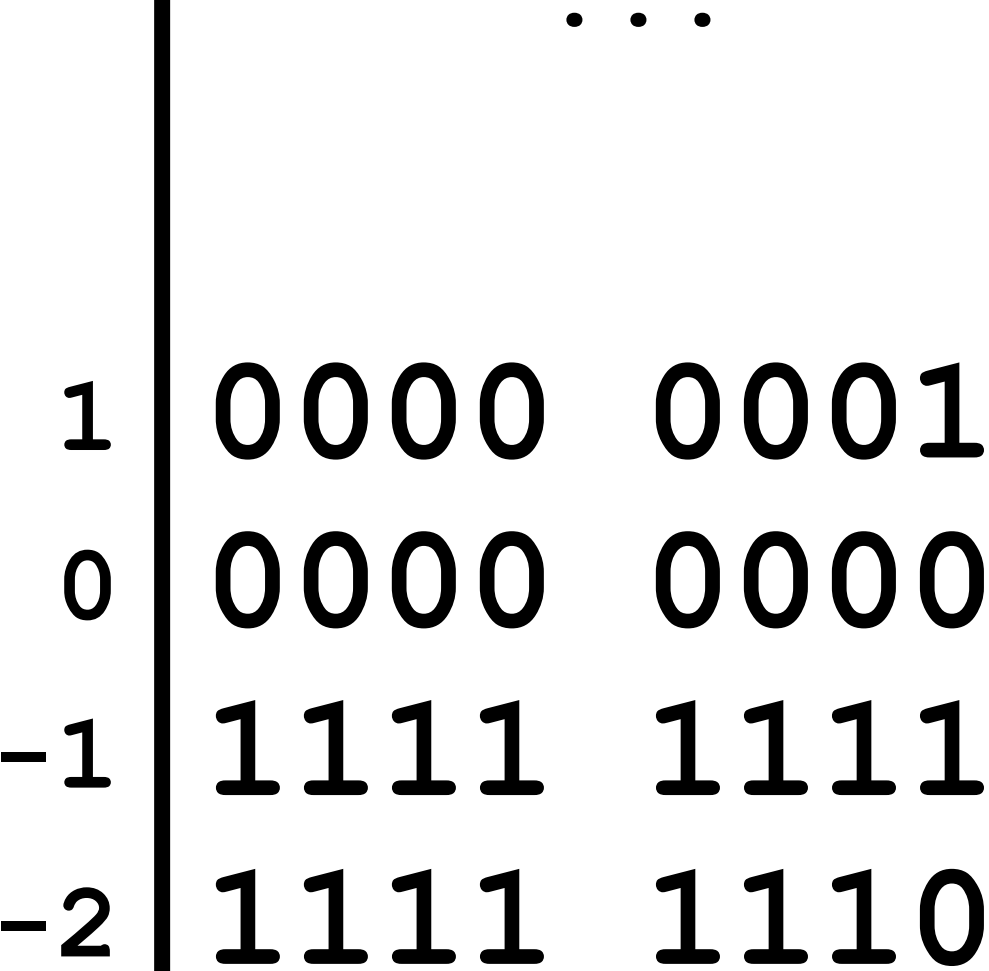
Two's complement: two ways



Negative Numbers

Two's complement: range

INT_MAX



1	0000	0001
0	0000	0000
-1	1111	1111
-2	1111	1110

Negative Numbers

Two's complement: range

INT_MAX ↑ 0111 1111

...

1 0000 0001

0 0000 0000

-1 1111 1111

-2 1111 1110

...

Negative Numbers

Two's complement: range

INT_MAX ↑ 0111 1111

...

1 0000 0001

0 0000 0000

-1 1111 1111

-2 1111 1110

...

INT_MIN

Negative Numbers

Two's complement: range

INT_MAX ↑ 0111 1111

...

1 0000 0001

0 0000 0000

-1 1111 1111

-2 1111 1110

...

INT_MIN | 1000 0000

Negative Numbers

Two's complement: range

INT_MAX ↑ 0111 1111 <-- $2^7 - 1$

...

1 0000 0001

0 0000 0000

-1 1111 1111

-2 1111 1110

...

INT_MIN | 1000 0000 <-- -2^7

Signed Arithmetics

INT_MAX	↑	0111	1111
		...	
1		0000	0001
0		0000	0000
-1		1111	1111
-2		1111	1110
		...	
INT_MIN		1000	0000

	1111	1110
+	0000	0011

Signed Arithmetics

INT_MAX	↑	0111	1111
		...	
1		0000	0001
0		0000	0000
-1		1111	1111
-2		1111	1110
		...	
INT_MIN		1000	0000

11111110

+ 00000011

-23

Signed Arithmetics

INT_MAX	↑	0111	1111
		...	
1		0000	0001
0		0000	0000
-1		1111	1111
-2		1111	1110
		...	
INT_MIN		1000	0000

	1111	1110	
+	0000	0011	-2
			3
		1	

Signed Arithmetics

INT_MAX	↑	0111	1111
		...	
1		0000	0001
0		0000	0000
-1		1111	1111
-2		1111	1110
		...	
INT_MIN		1000	0000

	1111	1110	
+	0000	0011	
		01	

-2
3

Signed Arithmetics

INT_MAX	↑	0111	1111
		...	
1		0000	0001
0		0000	0000
-1		1111	1111
-2		1111	1110
		...	
INT_MIN		1000	0000

	1111	1110	
+	0000	0011	
		001	

-2
3

Signed Arithmetics

INT_MAX	↑	0111	1111
		...	
1		0000	0001
0		0000	0000
-1		1111	1111
-2		1111	1110
		...	
INT_MIN		1000	0000

11111110

+ 00000011

0001

1

1

1

-2

3

Signed Arithmetics

INT_MAX	↑	0111	1111
		...	
1		0000	0001
0		0000	0000
-1		1111	1111
-2		1111	1110
		...	
INT_MIN		1000	0000

11111110

+ 00000011

00001

-2

3

Signed Arithmetics

INT_MAX	↑	0111	1111
		...	
1		0000	0001
0		0000	0000
-1		1111	1111
-2		1111	1110
		...	
INT_MIN		1000	0000

	1111	1110	
+	0000	0011	
	1	1	1
	1	1	
	00	0001	

-2
3

Signed Arithmetics

INT_MAX	↑	0111	1111
		...	
1		0000	0001
0		0000	0000
-1		1111	1111
-2		1111	1110
		...	
INT_MIN		1000	0000

	1111	1110
+	0000	0011
	1 1 1 1	1 1
	000	0001

-2
3

Signed Arithmetics

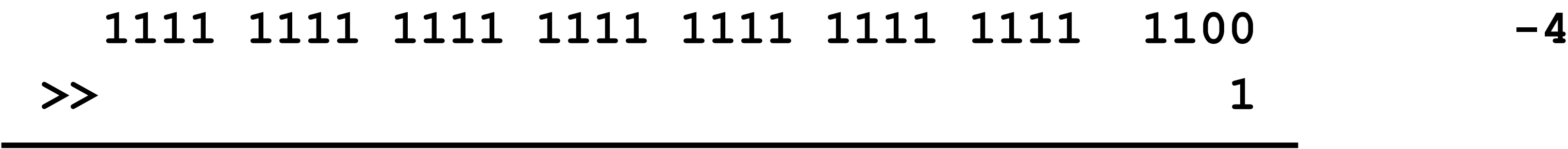
INT_MAX	↑	0111	1111
		...	
1		0000	0001
0		0000	0000
-1		1111	1111
-2		1111	1110
		...	
INT_MIN		1000	0000

$$\begin{array}{r}
 1111 \ 1110 \quad -2 \\
 + 0000 \ 0011 \quad 3 \\
 \hline
 0000 \ 0001
 \end{array}$$

- The hardware does not need to know/care whether the arithmetics is signed or unsigned
- $\text{INT_MAX} + 1 == \text{INT_MIN}$

Signed Arithmetics

Unsigned >>



Signed Arithmetics

Unsigned >>

1111	1111	1111	1111	1111	1111	1111	1100	-4
>>							1	
0111	1111	1111	1111	1111	1111	1111	1110	2147483646

- If we just add 0 in right shift, shifting a negative number results in a positive number.
- Ideally, we would like to keep the sign / shifting is dividing by 2.
- We repeat the most significant bit when doing a *signed* right shift.

Signed Arithmetics

Signed >>

>>

1111 1111 1111 1111 1111 1111 1111 1100

1

1111 1111 1111 1111 1111 1111 1111 1110

-4

Signed >>

	1111	1111	1111	1111	1111	1111	1111	1100	-4
>>								1	
	1111	1111	1111	1111	1111	1111	1111	1110	-2
	0000	0000	0000	0000	0000	0000	0000	0100	4
>>								1	

Signed Arithmetics

Signed >>

	1111	1111	1111	1111	1111	1111	1111	1100	-4
>>								1	
	1	111	1111	1111	1111	1111	1111	1110	-2
	0000	0000	0000	0000	0000	0000	0000	0100	4
>>								1	
	0	000	0000	0000	0000	0000	0000	0010	2

- If the first bit is 0, right shifting fills with 0.
- If the first bit is 1, right shifting fills with 1.

Signed Arithmetics

Signed >>

- In C, *signed* right shift and *unsigned* right shift looks exactly the same
- C will use *signed* shift if the variable being shifted is a *signed* integer

```
int x = ~0;  
x = x >> 4;  
printf("0x%x\n", x);
```

0xffffffff

```
unsigned int x = ~0;  
x = x >> 4;  
printf("0x%x\n", x);
```

0x0fffffff

Signed Arithmetics

- Signed extension works the same way:

```
short x = ~0;  
int y = x;  
printf("0x%x\n", y);
```

0xffffffff

```
unsigned short x = ~0;  
unsigned int y = x;  
printf("0x%x\n", y);
```

0x0000ffff

Machine Arithmetics

Summary

- For a 32-bit integer:
 - Range: unsigned $[0, 2^{32} - 1]$; signed $[-2^{31}, 2^{31} - 1]$.
 - Meaning of the highest bit: unsigned 2^{31} , signed -2^{31} .
- Negate a number: flips all bits and add 1 ($-n == \sim n + 1$).
- Most arithmetic operators work regardless of the signedness of the integer except right shift:
 - Signed right shift (arithmetic right shift): fills with the highest bit.
 - Unsigned right shift (logical right shift): fills with 0.
 - Left shift is the same.

Octal Number



Octal Number

- Base-8 numbers, 0-7
- Corresponds to 3 bits
- In C, prefix by 0, e.g. 0123, is 83
- Not very common

Bits

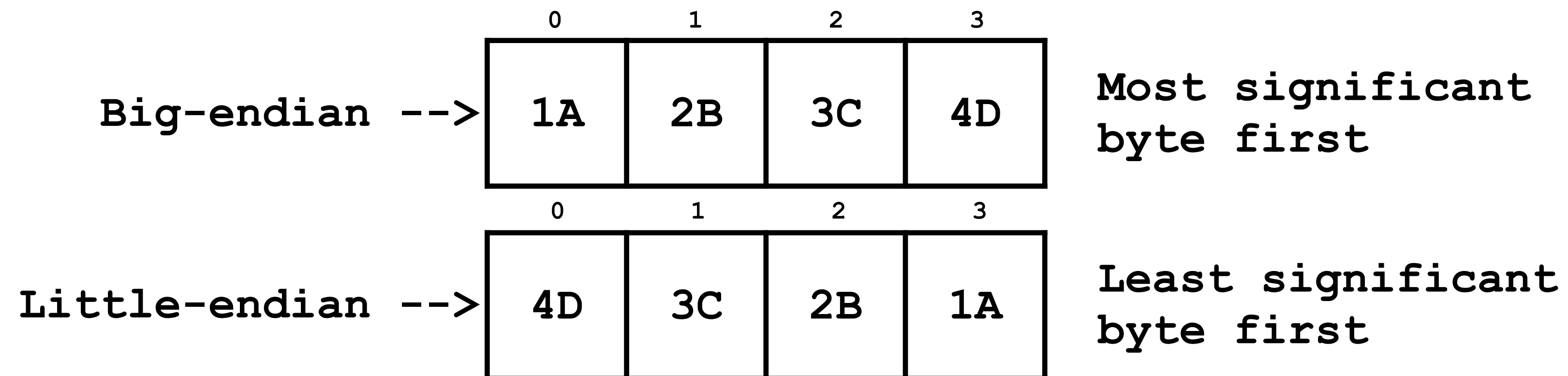
- A bit answers a yes/no question
 - Bits have no inherent meanings; to read encoding, need to know the intended interpretation
 - Any device that can be one of two states can encode a bit.
- Binary and hexadecimal numbers
 - $xyz_m = x \cdot m^2 + y \cdot m^1 + z \cdot m^0$
 - One hex digit corresponds to 4 binary digits (bits) -- easy conversion between the two
- Positive and negative numbers
- Arithmetics
 - $+$, $-$, $<<$, $\&$, $|$, $\wedge$, $\sim$, two $>>$, negate
 - Bit-packing

Enum

Demo

Endian

- We think of an integer as one atomic value:
 - `int x = 0x1A2B3C4D;`
- But if an integer has 4 bytes and each byte is addressable, which of the 4 bytes is stored first?



Endian

- Is my machine little-endian or big-endian?
- Let's find out!

Endian

- Our machine is little-endian?????
- We usually write numbers in big-endian: 345 is three hundred and forty-five
- But there are some advantages for little-endian:
 - comparing two numbers of different length (long and int e.g.)
 - 4E3C2B1A
 - 4E3C2B1A00000000
 - addition, subtraction circuits work from low to high
 - etc.

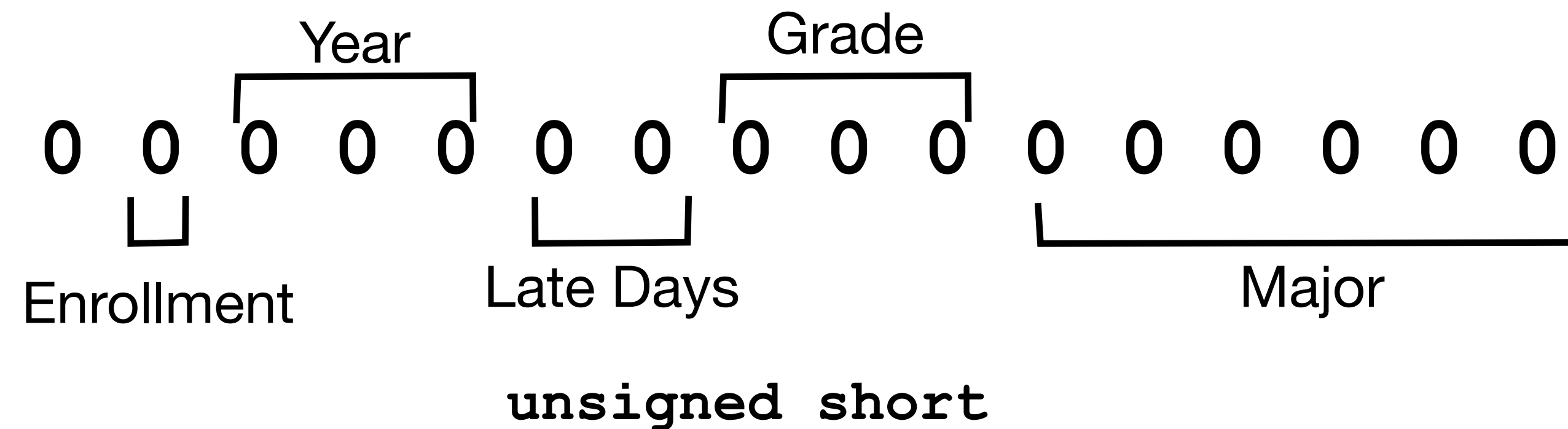
Endian

Does it matter?

- Mostly we don't care. Unless you do memory trickery, variables work as you would expect
- However, when we serialize data into byte sequences, you need to pay extra attention:
 - Writing a number to a file
 - Sending a number over a network
- You and the reader must agree on byte order
 - For this purpose, *network byte order* is defined for TCP/IP

Choices

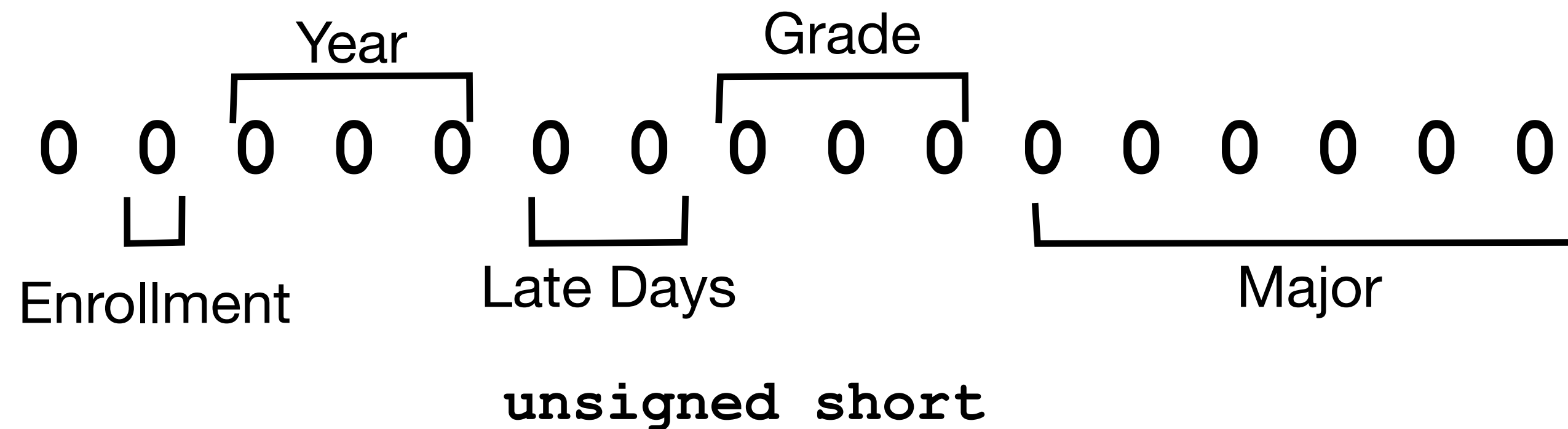
Interlude



- How do we assign bit patterns to majors?
- Anthropology - 0, Architectural Studies - 1, Art History - 2, ...
- One has to look up on a table to figure out which major this is :(

Choices

Interlude: enum

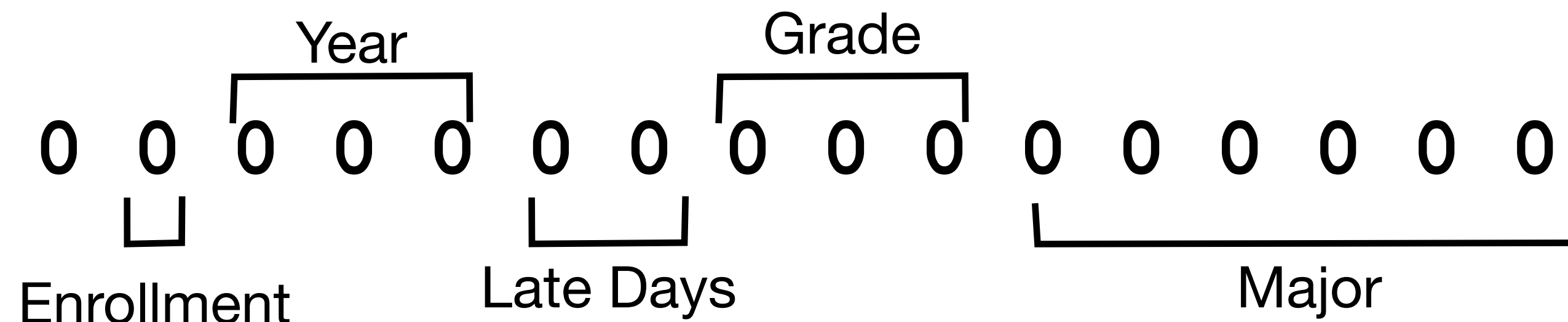


```
enum major {  
    ANTHROPOLOGY,  
    ARCHITECTURAL_STUDIES,  
    ART_HISTORY,  
    ASTRONOMY_ASTROPHYSICS,  
    BIG_PROBLEMS,  
    BIOLOGICAL_CHEMISTRY,  
    ...  
};
```

- Nothing fancy here: C just assigns an integer sequentially for each choice.
- Can use them as global constants
- **if** (major == 2) { ... }
- **if** (major == ART_HISTORY) { ... }

Choices

Interlude: enum



unsigned short

```
enum major {  
    ANTHROPOLOGY,  
    ARCHITECTURAL_STUDIES,  
    ART_HISTORY,  
    ASTRONOMY_ASTROPHYSICS,  
    BIG_PROBLEMS,  
    BIOLOGICAL_CHEMISTRY,  
    ...  
};
```

```
enum major student_major = STATISTICS;
```

```
enum major student_major = 3;
```

clang **will not**
complaint about this.
But this is bad style.

Choices

Interlude: switch

```
if (major == ANTHROPOLOGY) {  
    ...  
} else if (major == ARCHITECTURAL_STUDIES) {  
    ...  
} else if (major == ART_HISTORY) {  
    ...  
} else if (major == ASTRONOMY_ASTROPHYSICS) {  
    ...  
} else if (major == BIG_PROBLEMS) {  
    ...  
} ...
```

```
switch (major) {  
case ANTHROPOLOGY:  
    ...;  
    break;  
case ARCHITECTURAL_STUDIES:  
    ...;  
    break;  
case ART_HISTORY:  
    ...;  
    break;  
case BIOLOGICAL_CHEMISTRY:  
    ...;  
    break;  
case BIG_PROBLEMS:  
    ...;  
    break;  
default:  
    break;  
}
```

break signals the end of a case. Without **break**, C will execute the next case, *falling through* another case.

The default branch is run when the major matches none of the above cases.

Choices

Interlude: `switch`

- Why `switch`?
- Cleaner code
- More efficient than `if ... else if ...` chain:
 - C stores the branches in a table, and `switch` will jump to the branch instead of comparing one by one
 - `switch(x)`, `x` has to be an integer.
- `break` is critical! Forgetting the `break` is really difficult to debug.

Choices

Interlude: switch

- Demo! if we have time