

Why Standard ML?

A language particularly suited to compiler implementation.

- Efficiency
- Safety
- Simplicity
- Higher-order functions
- Static type checking with type inference
- Polymorphism
- Algebraic types and pattern matching
- Modularity
- Garbage collection
- Exception handling
- Libraries and tools

Using the SML/NJ Compiler

- *Type “sml” to run the SML/NJ compiler*

Normally installed in /usr/local/bin, which should be in your PATH.

- *Cntl-d exits the compiler, Cntl-c interrupts execution.*

- *Three ways to run ML programs:*

1. type in code in the interactive read-eval-print loop

2. edit ML code in a file, say `foo.sml`, then type command

```
use "foo.sml";
```

3. use Compilation Manager (CM):

```
CM.make "sources.cm";
```

ML Tutorial I

Expressions

- *Integers:* `3`, `54`, `~3`, `~54`
- *Reals:* `3.0`, `3.14159`, `~3.2E2`
- *Overloaded arithmetic operators:* `+`, `-`, `*`, `/`, `<`, `<=`
- *Booleans:* `true`, `false`, `not`, `orelse`, `andalso`
- *Strings:* `"abc"`, `"hello world\n"`, `x^".sml"`
- *Lists:* `[]`, `[1,2,3]`, `["x","str"]`, `1::2::nil`
- *Tuples:* `()`, `(1,true)`, `(3,"abc",true)`
- *Records:* `{a=1,b=true}`, `{name="fred",age=21}`
- *conditionals, function applications, let expressions, functions*

ML Tutorial 2

Declarations: *binding a name to a value*

value bindings

```
val x = 3  
val y = x + 1
```

function bindings

```
fun fact n =  
    if n = 0 then 1  
    else n * fact(n-1)
```

Let expressions: *local definitions*

```
let decl in expr end
```

```
let val x = 3  
    fun f y = (y, x*y)  
    in f(4+x)  
end
```

ML Tutorial 3

Function expressions

The expression “**fn** *var* => *exp*” denotes a function with formal parameter *var* and body *exp*.

```
val inc = fn x => x + 1
```

is equivalent to

```
fun inc x = x + 1
```

ML Tutorial 4

Compound values

Tuples: $(\text{exp}_1, \dots, \text{exp}_n)$

```
(3, 4.5)
```

```
val x = ("foo", x*1.5, true)
```

```
val first = #1(x)
```

```
val third = #3(x)
```

Records: $\{\text{lab}_1 = \text{exp}_1, \dots, \text{lab}_n = \text{exp}_n\}$

```
val car = {make = "Ford", year = 1910}
```

```
val mk = #make car
```

```
val yr = #year car
```

ML Tutorial 5

Patterns

a form to decompose compound values, commonly used in value bindings and function arguments

val *pat* = *exp*

fun *f*(*pat*) = *exp*

variable patterns:

val *x* = 3

$\Rightarrow$ *x* = 3

fun *f*(*x*) = *x*+2

tuple and record patterns:

val *pair* = (3,4.0)

val (*x*,*y*) = *pair*

$\Rightarrow$ *x* = 3, *y* = 4.0

val {*make*=*mk*, *year*=*yr*} = *car*

$\Rightarrow$ *mk* = "Ford", *yr* = 1910

ML Tutorial 6

Patterns

wildcard pattern: `_` (*underscore*)

constant patterns: `3`, `"a"`

```
fun iszero(0) = true
  | iszero(_) = false
```

constructor patterns:

```
val list = [1,2,3]
val fst::rest = list
⇒ fst = 1, rest = [2,3]
```

```
val [x,_,y] = list
⇒ x = 1, y = 3
```


ML Tutorial 7

Pattern matching

match rule: $pat \Rightarrow exp$

$$\text{match:} \quad pat_1 \Rightarrow exp_1 \mid \dots \mid pat_n \Rightarrow exp_n$$

When a match is applied to a value v , we try rules from left to right, looking for the first rule whose pattern matches v . We then bind the variables in the pattern and evaluate the expression.

case expression: **case** *exp* **of** *match*

function expression: **fn** *match*

```

clausal functional defn:   fun f pat1 = exp1
                           | f pat2 = exp2
                           | ...
                           | f pat2 = exp2

```

ML Tutorial 8

Pattern matching examples (function definitions)

```
fun length l = (case l
  of [] => 0
    | [a] => 1
    | _ :: r => 1 + length r
  (* end case *))
```

```
fun length [] = 0
  | length [a] = 1
  | length (_ :: r) = 1 + length r
```

```
fun even 0 = true
  | even n = odd(n-1)
```

```
and odd 0 = false
  | odd n = even(n-1)
```

ML Tutorial 9

Types

basic types: `int, real, string, bool`
 `3 : int,    true : bool,    "abc" : string`

function types: `$t_1 \rightarrow t_2$`
 `even: int -> bool`

product types: `$t_1 * t_2$  , unit`
 `(3,true): int * bool,    (): unit`

record types: `{lab1 : t1, ... , labn : tn}`
 `car: {make : string, year : int}`

type operators: `t list    (for example)`
 `[1,2,3] : int list`

ML Tutorial 10

Type abbreviations

```
type tycon = ty
```

examples:

```
type point = real * real
```

```
type line = point * point
```

```
type car = {make: string, year: int}
```

```
type tyvar tycon = ty
```

examples:

```
type 'a pair = 'a * 'a
```

```
type point = real pair
```

ML Tutorial I I

Datatypes

datatype $tycon = con_1 \text{ of } ty_1 \mid \dots \mid con_n \text{ of } ty_n$

This is a *tagged union* of variant types ty_1 through ty_n . The tags are the *data constructors* con_1 through con_n .

The data constructors can be used both in expressions to build values, and in patterns to deconstruct values and discriminate variants.

The “**of** ty ” can be omitted, giving a nullary constructor.

Datatypes can be *recursive*.

datatype $intlist = Nil \mid Cons \text{ of } int * intlist$

ML Tutorial 12

Datatype example

```
datatype btree = LEAF
                | NODE of int * btree * btree

fun depth LEAF = 0
    | depth (NODE(_,t1,t2)) =
        max(depth t1, depth t2) + 1

fun insert(LEAF,k) = NODE(k,LEAF,LEAF)
    | insert(NODE(i,t1,t2),k) =
        if k > i then NODE(i,t1,insert(t2,k))
        else if k < i then NODE(i,insert(t1,k),t2)
        else NODE(i,t1,t2)

(* in-order traversal of btrees *)
fun inord LEAF = []
    | inord(NODE(i,t1,t2)) =
        inord(t1) @ (i :: inord(t2))
```

ML Tutorial 13

Representing programs as datatypes

```
type id = string
```

```
datatype binop = PLUS | MINUS | TIMES | DIV
```

```
datatype stm = SEQ of stm * stm  
              | ASSIGN of id * exp  
              | PRINT of exp list
```

```
and exp = VAR of id  
         | CONST of int  
         | BINOP of binop * exp * exp  
         | ESEQ of stm * exp
```

```
val prog =  
    SEQ(ASSIGN("a", BINOP(PLUS, CONST 5, CONST 3)),  
        PRINT[VAR "a"])
```

ML Tutorial 14

Computing properties of programs: size

```
fun sizeS (SEQ(s1,s2)) = sizeS s1 + sizeS s2
  | sizeS (ASSIGN(i,e)) = 2 + sizeE e
  | sizeS (PRINT es) = 1 + sizeEL es

and sizeE (BINOP(_,e1,e2)) = sizeE e1 + sizeE e2 + 2
  | sizeE (ESEQ(s,e)) = sizeS s + sizeE e
  | sizeE _ = 1

and sizeEL [] = 0
  | sizeEL (e::es) = sizeE e + sizeEL es

sizeS prog  $\Rightarrow$  8
```


Types Review

Primitive types

unit, int, real, char, string, ..., instream, outstream, ...

Composite types

unit, tuples, records

function types

Datatypes

types and n-ary type operators, tagged unions, recursive

nominal type equality

bool, list

user defined: trees, expressions, etc.

Type Abbreviations

types and n-ary type operators

structural type equality

type 'a pair = 'a * 'a

Type Inference

When defining values (including functions), types do not need to be declared – they will be inferred by the compiler.

```
- fun f x = x + 1;  
  val f = fn : int -> int
```

Inconsistencies will be detected as type errors.

```
- if 1<2 then 3 else 4.0;  
stdIn:2.1-2.23 Error: types of rules don't agree  
  earlier rule(s): bool -> int  
  this rule: bool -> real  
  in rule:  
    false => 4.0
```

Type Inference

In some cases involving record field selections, explicit type annotations (called ascriptions) may be required

```
- datatype king = {name: string,  
                  born: int,  
                  crowned: int,  
                  died: int,  
                  country: string}
```

```
- fun lifetime(k: king) =  
  =      #died k - #born k;  
val lifetime = fn : king -> int
```

```
- fun lifetime({born,died,...}: king) =  
  =      died - born;  
val lifetime = fn : king -> int
```

*partial record
pattern*



Polymorphic Types

The most general type is inferred, which may be polymorphic

```
- fun ident x = x;  
val ident = fn : 'a -> 'a
```

```
- fun pair x = (x, x);  
val ident = fn : 'a -> 'a * 'a
```

```
- fun fst (x, y) = x;  
val ident = fn : 'a * 'b -> 'a
```

```
- val foo = pair 4.0;  
val foo : real * real
```

```
- fst foo;  
val it = 4.0: real
```

Polymorphic Types

The most general type is inferred, which may be polymorphic

```
- fun ident x = x;
```

```
val ident = fn : 'a -> 'a
```

type variable



```
- fun pair x = (x, x);
```

```
val ident = fn : 'a -> 'a * 'a
```

polymorphic type



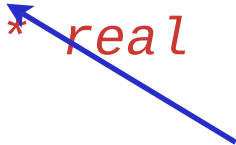
```
- fun fst (x, y) = x;
```

```
val ident = fn : 'a * 'b -> 'a
```

```
- val foo = pair 4.0;
```

```
val foo : real * real
```

: real -> real * real



```
- fst foo;
```

```
val it = 4.0: real
```

Polymorphic Data Structures

```
- infixr 5 ::  
- datatype 'a list = nil  
                  | :: of 'a * 'a list  
  
- fun hd nil = raise Empty  
=    | hd (x::_) = x;  
val hd = fn : 'a list -> 'a  
  
- fun length nil = 0  
=    | length (_::xs) = 1 + length xs;  
val length = fn : 'a list -> int  
  
- fun map f nil = nil  
=    | map f (x::xs) = f x :: map f xs;  
val map = fn : ('a -> 'b) -> 'a list -> 'b list
```

More Pattern Matching

Layered Patterns: x as pat

```
( * merging two sorted lists of ints * )  
fun merge(x, nil) = x  
  | merge(nil, y) = y  
  | merge(l as x::xs, m as y::ys) =  
    if x < y then x :: merge(xs,m)  
    else if y < x then y :: merge(l,m)  
    else x :: merge(xs,ys);  
val merge = fn : int list * int list -> int list
```

Note: although < is overloaded, this definition is unambiguously typed with the lists assumed to be int lists because the < operator defaults to the int version (of type `int*int->bool`).

Exceptions

```
- 5 div 0;                                (* primitive failure *)
```

```
uncaught exception Div
```

```
exception NotFound of string;    (* control structure *)
```

```
type 'a dict = (string * 'a) list
```

```
fun lookup (s,nil) = raise (NotFound s)
```

```
  | lookup (s,(a,b)::rest) =
```

```
    if s = a then b else lookup (s,rest)
```

```
val lookup: string * 'a dict -> 'a
```

```
val dict = [("foo",2), ("bar",~1)];
```

```
val dict: string * int list           (* == int dict *)
```

```
val x = lookup("foo",dict);
```

```
val x = 2 : int
```

```
val y = lookup("moo",dict);
```

```
uncaught exception NotFound
```

```
val z = lookup("moo",dict) handle NotFound s =>
```

```
  (print ("can't find "^s^"\n"); 0)
```

```
can't find moo
```

```
val z = 0 : int
```


References and Assignment

```
type 'a ref
val ref : 'a -> 'a ref
val ! : 'a ref -> 'a
val := : 'a ref * 'a -> unit
```

```
val linenum = ref 0;    (* create updatable ref cell *)
val linenum = ref 0 : int ref
```

```
fun newLine () = linenum := !linenum + 1;  (* increment it *)
val newline = fn : unit -> unit
```

```
fun lineCount () = !linenum;  (* access ref cell *)
val lineCount = fn : unit -> int
```

```
local val x = 1
  in fun new1 () = let val x = x + 1 in x end
end  (* new1 always returns 2 *)
```

```
local val x = ref 1
  in fun new2 () = (x := !x + 1; !x)
end  (* new2 returns 2, 3, 4, ... on successive calls *)
```

Simple Modules -- *Structure*

```
structure Ford =  
struct  
  type car = {make: string, built: int}  
  val first = {make = "Ford", built: 1904}  
  fun mutate ({make,built}: car) year =  
    {make = make, built = year}  
  fun built ({built,...}: car) = built  
  fun show (c) = if built c < built first then " - "  
    else "(generic Ford)"  
end
```

```
structure Year =  
struct  
  type year = int  
  val first = 1900  
  val second = 2000  
  fun newYear(y: year) = y+1  
  fun show(y: year) = Int.toString y  
end
```

```
structure MutableCar =  
struct  
  structure C = Ford  
  structure Y = Year  
end
```

*A structure is an
encapsulated, named,
collection of declarations*

Module Interfaces -- *Signature*

```
signature MANUFACTURER =  
sig  
  type car  
  val first : car  
  val built : car -> int  
  val mutate : car -> int -> car  
  val show : car -> string  
end
```

```
signature YEAR =  
sig  
  eqtype year  
  val first : year  
  val second : year  
  val newYear : year -> year  
  val show : year -> string  
end
```

```
signature MCSIG =  
sig  
  structure C : MANUFACTURER  
  structure Y : YEAR  
end
```

*A signature is a
collection of specifications
for module components --
types, values, structures*

Signature Matching

```
structure Year1 : YEAR =  
struct  
  type year = int  
  type decade = string  
  val first = 1900  
  val second = 2000  
  fun newYear(y: year) = y+1  
  fun leap(y: year) = y mod 4 = 0  
  fun show(y: year) = Int.toString y  
end
```

Structure *S* matches SIG if *S*
if every spec in SIG is
matched by a component of *S*.

S can have more components
than are specified in SIG.

```
structure MCar : MCSIG = MutableCar
```

```
val classic = Year1.show 1968  
val antique = MCar.Y.show 1930
```

Use the *dot notation* to access
components of structures.

```
val x = Year1.leap(Year1.first)
```

Can't access components not
specified in signature.

Module Functions -- *Functors*

```
signature ORD =  
sig  
  type t  
  val less : t * t -> bool  
end
```

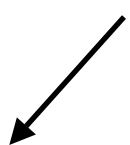
```
functor Sort(X: ORD) =  
struct  
  fun insert(x,nil) = [x]  
    | insert(x,l as y::ys) =  
      if X.less(x,y) then x::l  
      else y::insert(x,ys)  
  fun sort (m : X.t list) = foldl insert nil m  
end
```

Sort is a *parameterized module*, with parameter
X: ORD

```
structure IntOrd : ORD =  
struct  
  val t = int  
  val less = Int.<  
end
```

```
structure IntSort = Sort(IntOrd)
```

functor application



Input/Output

```
structure TextIO : sig
```

```
type instream          (* an input stream *)  
type ostream          (* an output stream *)
```

```
val stdin : instream    (* standard input *)  
val stdout : ostream    (* standard output *)  
val stderr : ostream    (* standard error *)
```

```
val openIn: string -> instream    (* open file for input *)  
val openOut: string -> ostream    (* open file for input *)  
val openAppend: string -> ostream (* open file for appending *)
```

```
val closeIn: instream -> unit      (* close input stream *)  
val closeOut: ostream -> unit     (* close output stream *)
```

```
val output: ostream * string -> unit (* output a string *)
```

```
val input: instream -> string      (* input a string *)  
val inputLine: instream -> string  (* input a line *)
```

```
.....
```

```
end
```

The Unix structure

```
signature UNIX =  
  sig  
    type proc  
    val execute : (string * string list) -> proc  
    val streamsOf : proc -> (TextIO.instream * TextIO.outstream)  
    val reap : proc -> Posix.Process.exit_status  
    ...  
  end  
  
structure Unix : UNIX = ...
```

Note: the Unix structure in SML/NJ has a different signature from that given in the SML'97 Basis documentation.

An example of using the Unix structure

```
structure U = Unix
```

```
fun date () = let  
    val pid = U.execute("/bin/date", [])  
    val (inS, outS) = U.streamsOf pid  
    val _ = TextIO.closeOut outS  
    val now = TextIO.inputLine inS  
in  
    TextIO.closeIn inS;  
    U.reap pid;  
    now  
end
```


Modules --- type abstraction

Consider the problem of providing *unique* identifiers.

```
signature UID =  
  sig  
    type uid  
    val same : (uid * uid) -> bool  
    val compare : (uid * uid) -> order  
    val gensym : unit -> uid  
  end
```

Modules --- type abstraction

```
structure UId :> UID =  
  struct  
    type uid = int  (* abstract *)  
    fun same (a : uid, b) = (a = b)  
    val compare = Int.compare  
  
    val count = ref 0  (* hidden *)  
    fun gensym () = let  
      val id = !count  
    in  
      count := id + 1;  
      id  
    end  
  
  end
```