

Basic Blocks and Traces

Lecture 8

Canonical Trees

```
signature CANON =
```

```
sig
```

```
  val linearize : Tree.stm -> Tree.stm list
```

```
  val basicBlocks : Tree.stm list ->  
    (Tree.stm list list * Tree.label)
```

```
  val traceSchedule : Tree.stm list list * Tree.label ->  
    Tree.stm list
```

```
end (* signature CANON *)
```

Canonical Trees

signature CANON =
sig

```
val linearize : Tree.stm -> Tree.stm list
  (* From an arbitrary Tree statement, produce a list of
     canonical trees satisfying the following properties:
     1. No SEQ's or ESEQ's
     2. The parent of every CALL is an EXP(..) or a
        MOVE(TEMP t,..)
  *)
```

```
val basicBlocks : Tree.stm list ->
  (Tree.stm list list * Tree.label)
```

```
val traceSchedule : Tree.stm list list * Tree.label ->
  Tree.stm list
```

```
end (* signature CANON *)
```

Basic Blocks

signature CANON =
sig

val linearize : Tree.stm -> Tree.stm list

val basicBlocks : Tree.stm list ->

(Tree.stm list list * Tree.label)

(* From a list of cleaned trees, produce a list of
basic blocks satisfying the following properties:

1. and 2. as above;

3. Every block begins with a LABEL;

4. A LABEL appears only at the beginning of a block;

5. Any JUMP or CJUMP is the last stm in a block;

6. Every block ends with a JUMP or CJUMP;

Also produce the "label" to which control will be passed
upon exit.

*)

val traceSchedule : Tree.stm list list * Tree.label ->
Tree.stm list

end (* signature CANON *)

Traces

signature CANON =
sig

val linearize : Tree.stm -> Tree.stm list

val basicBlocks : Tree.stm list ->
 (Tree.stm list list * Tree.label)

val traceSchedule : Tree.stm list list * Tree.label ->
 Tree.stm list

(* From a list of basic blocks satisfying properties 1-6,
along with an "exit" label, produce a list of stms such
that:
 1. and 2. as above;
 7. Every CJUMP(_,t,f) is immediately followed by LABEL f.
The blocks are reordered to satisfy property 7; also
in this reordering as many JUMP(T.NAME(lab)) statements
as possible are eliminated by falling through into
T.LABEL(lab).
*)

end (* signature CANON *)

Canonical Trees

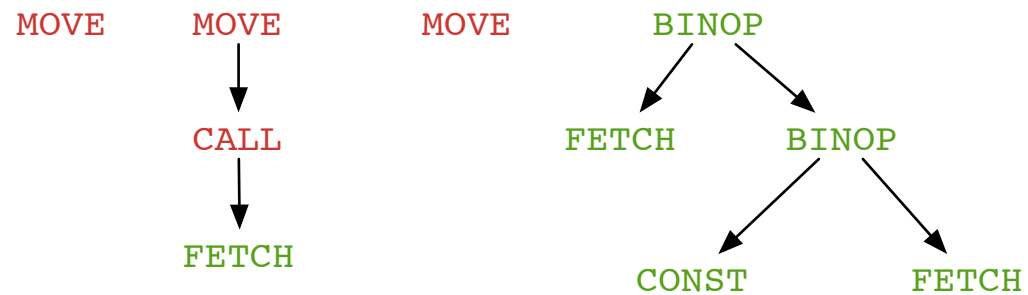
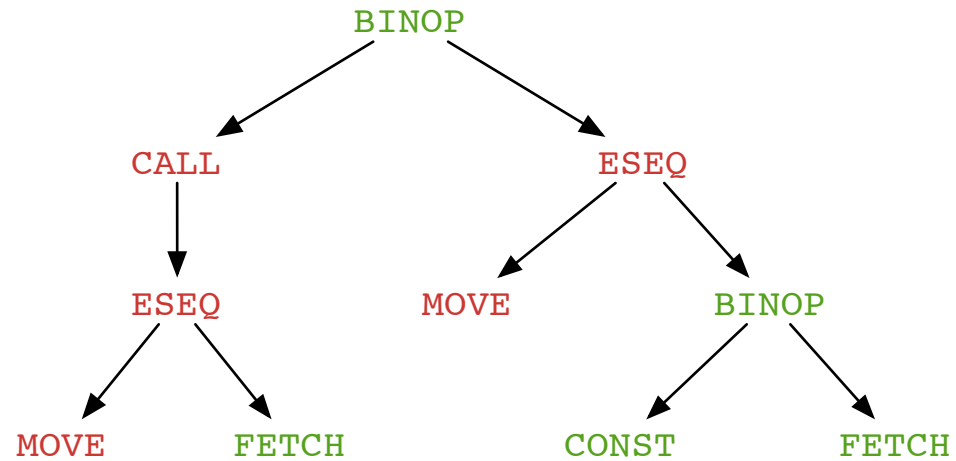
Canonical trees are those that:

1. Have no SEQ or ESEQ subterms
2. CALLs appear only as components of stms, not as subexpressions, i.e. a CALL node has parent of the form EXP(_) or MOVE(TEMP(t),_)

The idea is to separate out statements with side-effects from pure expressions. This allows freedom to change the order of evaluation in expressions and simplifies the interaction between expression evaluation (function calls in particular), and side-effects like assignment.

Linearization pulls stms and function calls to the top and front, linked with SEQ and ESEQ. Then the SEQ and ESEQ chain can be simplified to a list of canonical trees.

Canonical Tree Transformation



Canonical Tree Transforms

A number of term transformations are used to rearrange expressions into canonical form (Figure 8.1). E.g.:

$$\text{ESEQ}(s1, \text{ESEQ}(s2, e)) \implies \text{ESEQ}(\text{SEQ}(s1, s2), e)$$
$$\text{BINOP}(op, (\text{ESEQ}(s, e1), e2) \implies \\ \text{ESEQ}(s, (\text{BINOP}(op, e1, e2)))$$
$$\text{BINOP}(op, e1, (\text{ESEQ}(s, e2)) \implies \\ \text{ESEQ}(\text{MOVE}(\text{TEMP } t_{new}, e1), \\ \text{ESEQ}(s, (\text{BINOP}(op, \text{FETCH}(\text{TEMP } t_{new}), e2)))$$
$$\text{BINOP}(op, e1, (\text{ESEQ}(s, e2)) \implies \text{ESEQ}(s, \text{BINOP}(op, e1, e2))$$

if s and $e1$ commute (i.e. are noninterfering,
the effects performed by s will not change the
value computed by $e1$)