

Registers and Stack Frames

Lecture 6

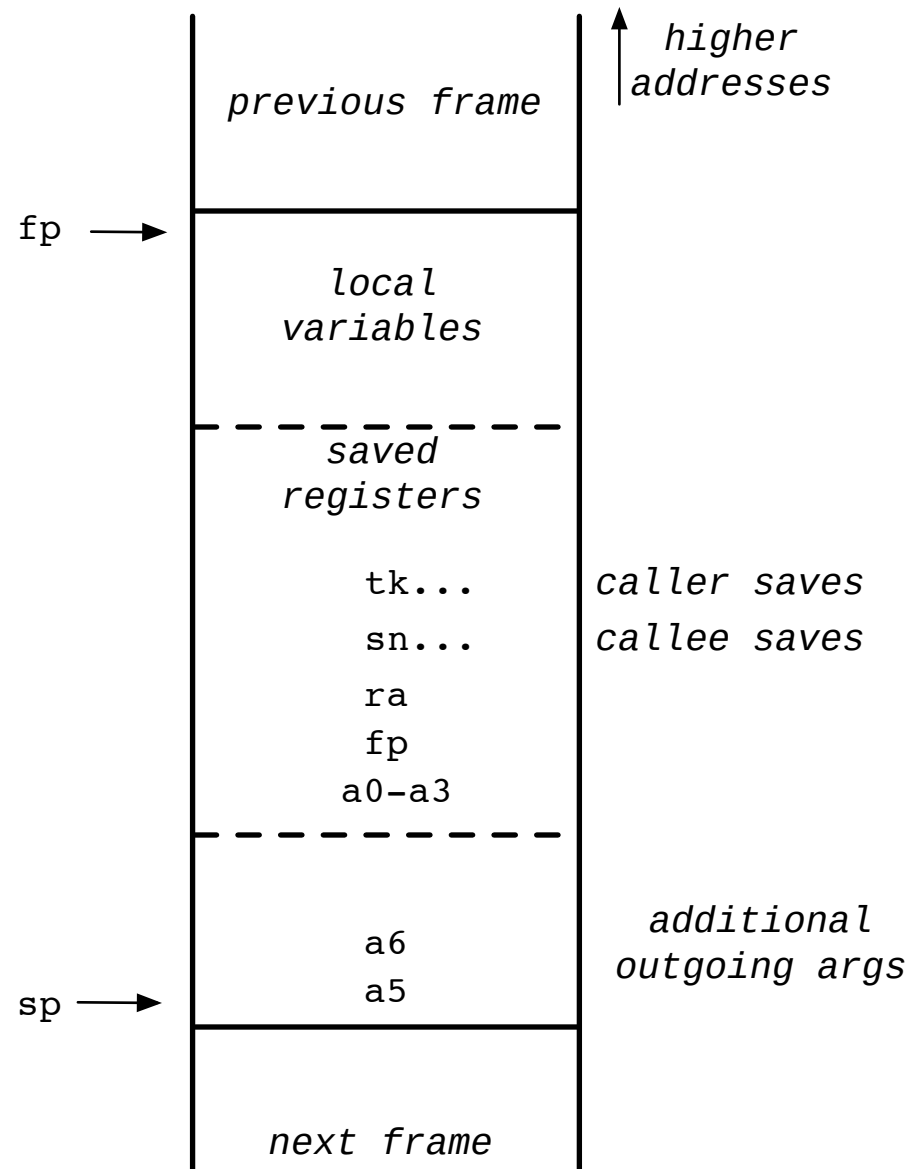
Function calls and stack frames

- *function invocation contexts are strictly nested*
 - stack mechanism is adequate for managing parameters and locals
- *Information needed for function calls*
 - arguments
 - in registers
 - in memory (via stack)
 - function return address
 - return values?
 - local variables of the function
 - temp space for saving registers

MIPS Registers

Hardware	Name	Description
\$0	zero	constant zero
\$1	at	assembler temporary
\$2-\$3	v0-v1	function return value
\$4-\$7	a0-a3	incoming args
\$8-\$15	t0-t7	temporaries
\$16-\$23	s0-s7	callee-saves temporaries
\$24-\$25	t8-t9	temporaries
\$26-\$27	k0-k1	exception handling
\$28	gp	global data pointer
\$29	sp	stack pointer
\$30	s8,fp	frame pointer
\$31	ra	return address

Stack Frame Layout



Assembly Example I

```
.text
.globl main

main:
    subu    $sp,$sp,24      # stack frame is 6 words
    sw      $ra,20($sp)     # save return address
    sw      $fp,16($sp)     # save frame pointer
    addu    $fp,$sp,20      # set new frame pointer

    li      $a0,6           # arg0: 6
    jal     f               # call f

    move    $a0,$v0         # arg0: return value from f
    jal     print_int       # call print_int

    lw      $ra,20($sp)     # restore return address
    lw      $fp,16($sp)     # restore frame pointer
    addu    $sp,$sp,24      # pop current frame
    j       $ra             # return to caller
```

Assembly Example 2

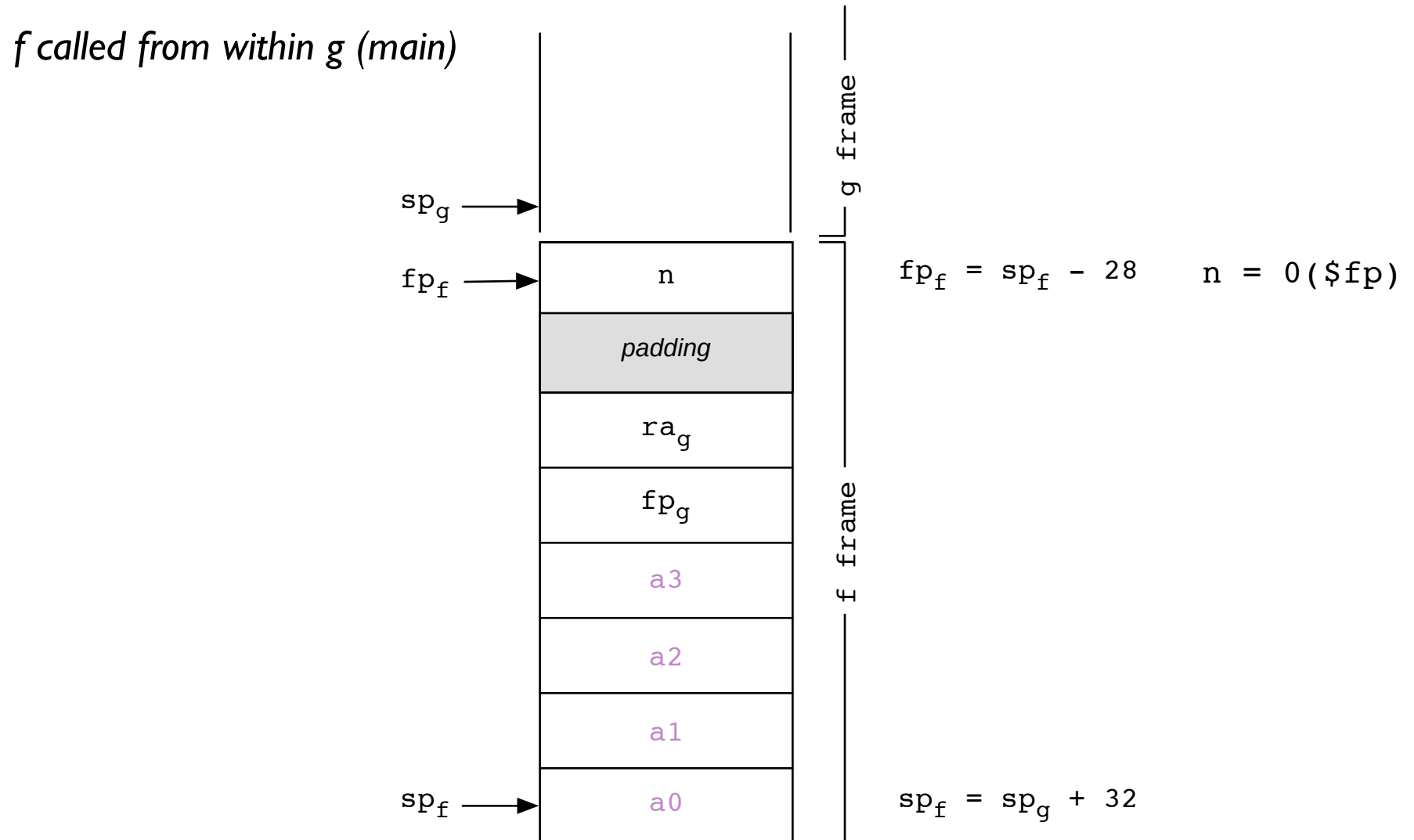
```
.text
f:
    subu    $sp,$sp,32      # allocate new frame (8 words)
    sw      $ra,20($sp)     # save return address
    sw      $fp,16($sp)     # save previous frame pointer
    addu    $fp,$sp,28      # define new frame pointer

    sw      $a0,0($fp)      # store n
    lw      $v0,0($fp)      # load n into v0
    bgtz    $v0,$L2         # branch to $L2 if n > 0
    li      $v0,1           # otherwise, return 1
    j       $L1             # jump to return code

$L2:    lw      $v1,0($fp)   # load n into v1
    subu    $a0,$v1,1       # a0 := v1 - 1 (= n - 1)
    jal     f               # call f recursively
    lw      $v1,0($fp)      # load n into v1
    mul     $v0,$v0,$v1     # v0 := f(n-1) * n

$L1:    lw      $ra,20($sp)  # restore return address
    lw      $fp,16($sp)     # restore frame pointer
    addu    $sp,$sp,32      # pop frame from stack
    j       $ra             # return to caller (main)
```

Example frame



Stack pointer must be aligned on two word boundary.

Escaping Variables

```
type tree = {key: string, left: tree, right: tree}

function prettyprint(tree: tree) : string =
  let
    var output := ""

    function write(s: string) =
      output := concat(output,s)

    function show(n: int, t: tree) =
      let function indent(s: string) =
        (for i := 1 to n
         do write(" "));
        output := concat(output,s); write("\n"))
      in if t=nil then indent(".")
        else (indent(t.key);
              show(n+1,t.left);
              show(n+1,t.left))
      end
  in show(0,tree);
  output
end
```

An *escaping variable* is a local variable of a function that occurs in the body of a nested function definition. E.g. **output**, **n**.

Computing Escaping Variables

- An escaping variable must be declared before the nested function that it appears in.
- The escaping property can be determined by comparing the function nesting depth of the variable's declaration with the nesting depths of its applied occurrences.
- The escaping property is recorded in the escape field (a `bool ref`) of the abstract syntax construct that declares it.
 - `VarDec` for variable declarations
 - `ForExp` for for loop index variables
 - `field` record for function parameters
- An environment can be used to record function nesting depth and escape field ref for each variable within its scope.

Static Links

- To access an escaping variable while executing the body of a nested function that accesses it, we will use a *static link*.
- A static link is the frame pointer of a frame (function activation record) for the latest call of the function statically enclosing the function currently being called.
- The static link will be passed as an added parameter, conventionally as the first parameter (a0).
- To access an escaping variable more than one nesting level deep, we need to follow a chain of static links.

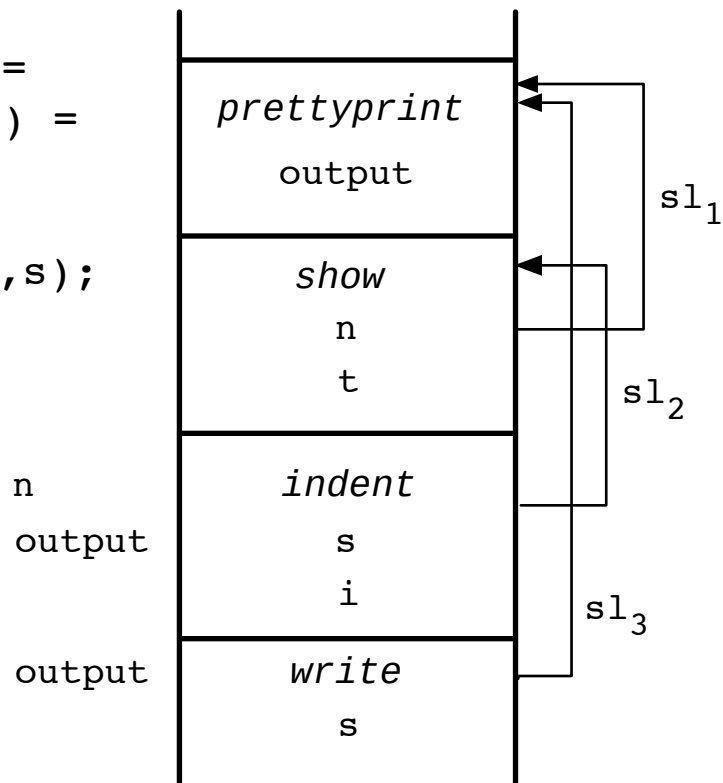
Static Links Example

```

function prettyprint(tree: tree) : string =
  let
    var output := ""

    function write(s: string) =
      output := concat(output,s)

    function show(n: int, t: tree) =
      let function indent(s: string) =
        (for i := 1 to n
         do write(" "));
        output := concat(output,s);
        write("\n")
      in if t=nil then indent(".")
        else (indent(t.key);
              show(n+1,t.left);
              show(n+1,t.left))
      end
    in show(0,tree);
    output
  end
end
  
```



Frames for Tiger

```
signature FRAME =  
sig
```

```
  type frame  
  type access
```

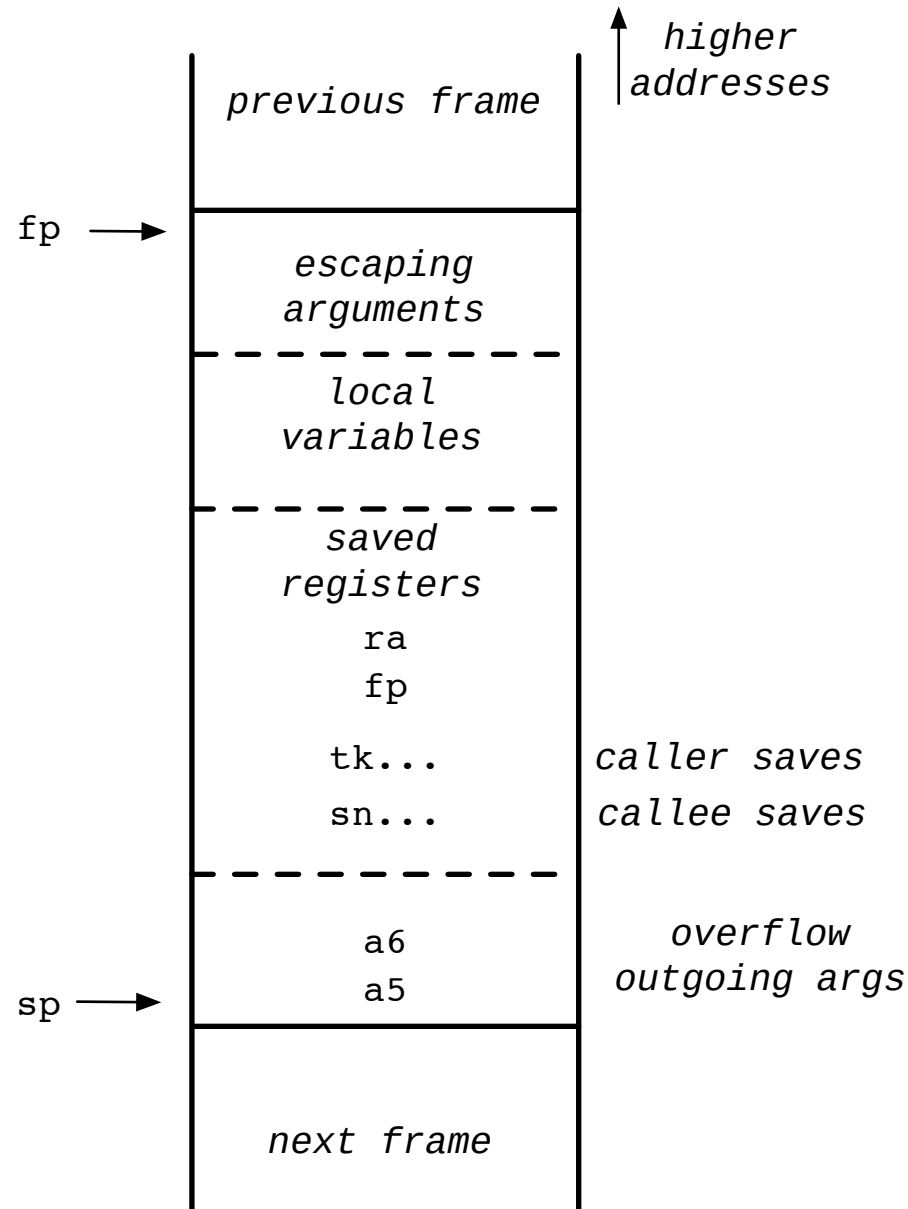
```
  val newFrame : {name: Temp.label,  (* label of called fun *)  
                  formals : (bool * string) list} -> frame  
    (* formals specified with an escape flag, and a description  
    * string *)
```

```
  val name : frame -> Temp.label  
    (* the function label of a frame *)  
  val formals : frame -> access list  
    (* access info for formal parameters *)  
  val allocLocal : frame -> (bool * string) -> access
```

```
  (* more to come ... *)
```

```
end (* signature FRAME *)
```

Frame Layout (2)



Incoming Arguments

```
function f(x0,x1,x2,x3,y4,y5) =  
  let function g(u: int): int =  
    x0 + x2 + u    /* x0 and x2 escaping */  
  in g(3)  
end
```

Arguments $x_0, \dots, x_3$ start in registers $\$a_0, \dots, \a_3 .
They are moved to temp registers or frame slots:

$x_0:\$a_0$	$\Rightarrow$	$0(\$fp)$	-- escaping
$x_1:\$a_1$	$\Rightarrow$	t_{17}	
$x_2:\$a_2$	$\Rightarrow$	$-4(\$fp)$	-- escaping
$x_3:\$a_3$	$\Rightarrow$	t_{18}	

Arguments y_4, y_5 are stored in previous frame:

$y_4:$	$4(\$fp)$
$y_5:$	$8(\$fp)$

