

Review and Preview

Lecture 12

Further Topics in Compilers

- *Advanced Language features*
 - Object-Oriented Languages
 - objects, classes
 - Functional Languages
 - function closures
 - lazy evaluation
 - advanced type systems
 - parametric polymorphism (ML, Haskell)
 - subtyping (O-O languages)
 - modules
 - continuations, coroutines, threads

Further Topics in Compilers

- *data flow and control flow analysis*
- *constant folding*
- *inlining functions*
- *loop optimization*
 - hoisting code out of loops, loop unrolling
- *instruction selection*
- *instruction scheduling*
- *array bounds checks*
- *delay slots, speculative execution*
- *instruction level parallelism (ILP), pipelining, function units*

Intermediate Languages

- *Static Single-Assignment form*
 - IR as functional language
- *Continuation-Passing Style (CPS)*
 - A-normal form
 - making important structure explicit
- *Typed intermediate languages*
 - FLINT, TILT, TAL, ...
- *Virtual machine systems (JVM, MSIL)*
 - JIT compilers

Runtime Systems

- *Memory management*
 - garbage collectors
 - mark-sweep (classical)
 - copying, generational, incremental, compacting
- *Traps and interrupts*
- *Input/Output*
- *OS services*
- *Threads*

Review

- *Lexical Analysis*
 - turning characters into tokens
 - regular expressions
 - NFAs and DFAs
 - lex (ml-lex)
 - r.e. patterns, actions
 - start states
 - handling strings with escapes
 - handling nested comments

Parsing

- *Context-free grammars*
- *Top-down parsers*
 - recursive descent
- *Bottom-up parsers*
 - LR grammars (LR(k), SLR, LALR(k))
 - constructing parser tables
- *Yacc (ml-yacc)*
 - terminals and nonterminals
 - grammar rules and actions
 - shift-reduce and reduce-reduce conflicts
 - precedence and associativity declarations

Abstract Syntax

- *Simple tree representation of logical structure*
 - types
 - expressions
 - statements
 - declarations
- *Naturally expressed using ML datatypes*
- *Basis for semantic (or static) analysis*

Type Checking

- *Types*
 - express structure of data
 - interface of functions
- *Type Environments*
 - map names (variables, functions) to their types
- *Typing rules*
 - relate expressions and types
- *Type checking*
 - check consistency
 - synthesize types of expressions

Semantic Analysis

- *Type checking*
- *Determining scope of names (types, variables, functions)*
- *Escape analysis*
 - global vs local variables
- *Forward references*
 - recursive types
 - recursive functions
 - two pass analysis (prebind names, then analyze definitions)

Intermediate Representation

- *Intermediate language or IR (tree.sig/sml)*
 - a lower-level tree representation of program structure
 - constructs similar to machine language
 - unlimited supply of temps, or abstract registers
 - temps and labels replace variables and functions
 - conditional and unconditional jumps express control flow constructs (if-then-else, while, for, break)
 - type *lexp* expresses memory accesses (l-values, r-values)
 - MOVE represents assignment, indexing, selection

Translation of Absyn to IR

- *translation environments*
 - maps names (variables and functions) to access info
- *separate expressions and statements*
 - gexp reunifies expressions, statements, and conditionals
 - coercions between different forms to satisfy context
- *recursive traversal of abstract syntax (similar to type checking, escape analysis)*
 - two pass treatment of recursive function declarations
 - types not involved
- *units of translation are “fragments”*
 - representing single function body (or top-level program)

Function Call Frames

- *Call Frames (aka Activation Records)*
 - store local information associated with a function call
 - arguments and local variables that “escape”
 - saved \$fp and \$ra registers
 - space for spilled temps and callee saves registers
 - space for excess outgoing arguments (beyond first 4)
- *frame record*
 - records information about function and its frame layout during compilation
 - manages allocation of slots for arguments, locals, spills
 - could store info on use of global variables and need for static link

Static Links

- *Need to compute access to nonlocal variables*
- *static link is frame pointer of frame of statically enclosing function*
- *passed to function as additional, implicit parameter*
 - not always needed
 - first argument, treated as escaping (found in $O(\$fp)$)
- *computed with the aid of “level” type*
 - chain of statically nested functions
 - translation environment maps function to its parent’s level

Basic Blocks & Trace Scheduling

- *Linearize code*
 - move statements (including calls) out of expressions
 - no side effects in expressions
 - flatten to statement list (possibly followed by final expression)
 - eliminates SEQ and ESEQ IR tree forms
- *Split into basic blocks*
 - straight-line code segments
 - start with label, end with (conditional or unconditional) jump
 - can be reordered without changing behavior
- *Trace scheduling*
 - sort basic blocks and concatenate them
 - arrange so jumps are followed by target labels when possible
 - arrange for false branch of cond. jump to follow jump

Liveness Analysis

- *live range*
 - a temp is live over a sequence of instructions between a definition and a use of that temp
- *liveness analysis determines the live ranges of temps*
 - calculates live-in and live-out sets of temps at each instruction
- *two temps **interfere** if they are both live at the same point*
 - they then have to coexist, and therefore can't occupy the same register
 - t1 and t2 interfere if t1 is defined at an instr and the other is live-out at that instr
- *interference graph records interference relation*
 - nodes represent temps, edges represent interference

Register Allocation

- *Assign registers (strings) to temps (allocation)*
- *Avoid assigning same register to temps that interfere*
 - color the interference graph using registers as colors
 - successively remove nodes of insignificant degree (*simplify*)
 - color them as they are restored
 - pick colors to maximize number of moves between temps of same color (these moves can then be eliminated)
 - If no nodes of insignificant degree
 - choose a node of minimal *spill cost* to *spill* (store in frame)
 - rewrite code to accomplish spilling
 - redo liveness analysis and interference graph coloring
- *Register coalescing*
 - attempt to coalesce nodes that are move related, if it doesn't make coloring harder

Final Exam

- *Open book, open notes, open code*
- *Wednesday, June 11, 10:30am-12:30pm*